

# Block I

## Apollo Guidance Computer (AGC)

How to build one in your basement

Part 3: Processing (PROC) Module

John Pultorak  
December, 2004

# Abstract

This report describes my successful project to build a working reproduction of the 1964 prototype for the Block I Apollo Guidance Computer. The AGC is the flight computer for the Apollo moon landings, and is the world's first integrated circuit computer.

I built it in my basement. It took me 4 years.

If you like, you can build one too. It will take you less time, and yours will be better than mine.

I documented my project in 9 separate .pdf files:

- Part 1            Overview: Introduces the project.
- Part 2            CTL Module: Design and construction of the control module.
- Part 3            PROC Module: Design and construction of the processing (CPU) module.
- Part 4            MEM Module: Design and construction of the memory module.
- Part 5            IO Module: Design and construction of the display/keyboard (DSKY) module.
- Part 6            Assembler: A cross-assembler for AGC software development.
- Part 7            C++ Simulator: A low-level simulator that runs assembled AGC code.
- Part 8            Flight Software: My translation of portions of the COLOSSUS 249 flight software.
- Part 9            Test & Checkout: A suite of test programs in AGC assembly language.

# Overview

The Processing Module (PROC) has 5 subsystems: PMI, ALU, CRG, INT, CTR

PMI (Processing Module external Interface)

The PMI interfaces other processing module subsystems to external AGC modules. 40-pin IDE connectors interface to the other CTL, MEM, and IO modules. Inputs from those modules are buffered to 1 LSTTL load.

ALU (Arithmetic Logic Unit)

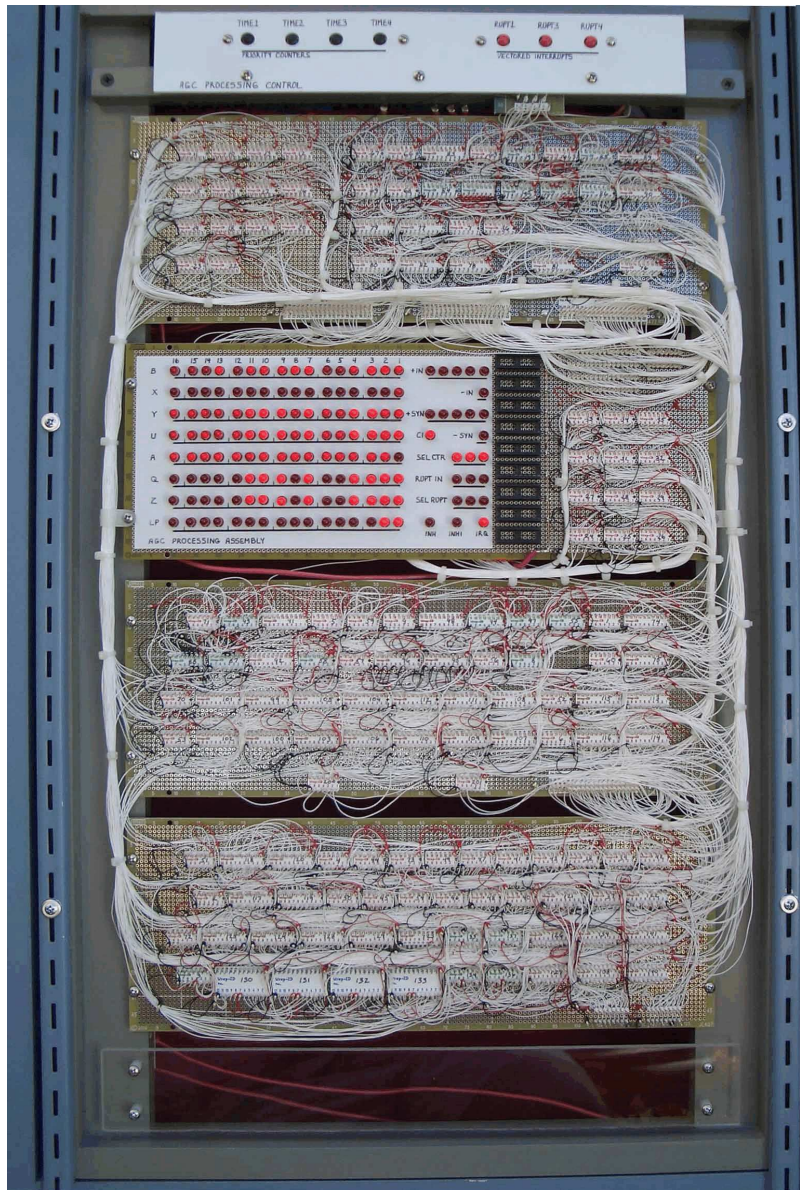
The ALU contains the 16-bit ADDER to perform 1's complement arithmetic and increment the program counter (Z register). The ALU also contains the B and C registers, and logic to inclusive OR the contents of the them with the ADDER. The inclusive OR occurs when the control module (CTL) issues signals to READ the contents of both registers onto the READ bus simultaneously. The ALU transfers data from the READ bus to the WRITE bus for register-to-register transfers and can force the data lines of the WRITE bus to specific states to gate various constants into the AGC registers

CRG (Central Register)

The AGC has four 16-bit "central registers" for general computational use. These are the accumulator (A) for general computation; the program counter (Z) which contains the address of the next instruction; the Q register holding the remainder in the DV instruction, and the return address after TC instructions; and the LP register used to hold the lower product after MP instructions.

INT (Interrupt Priority)

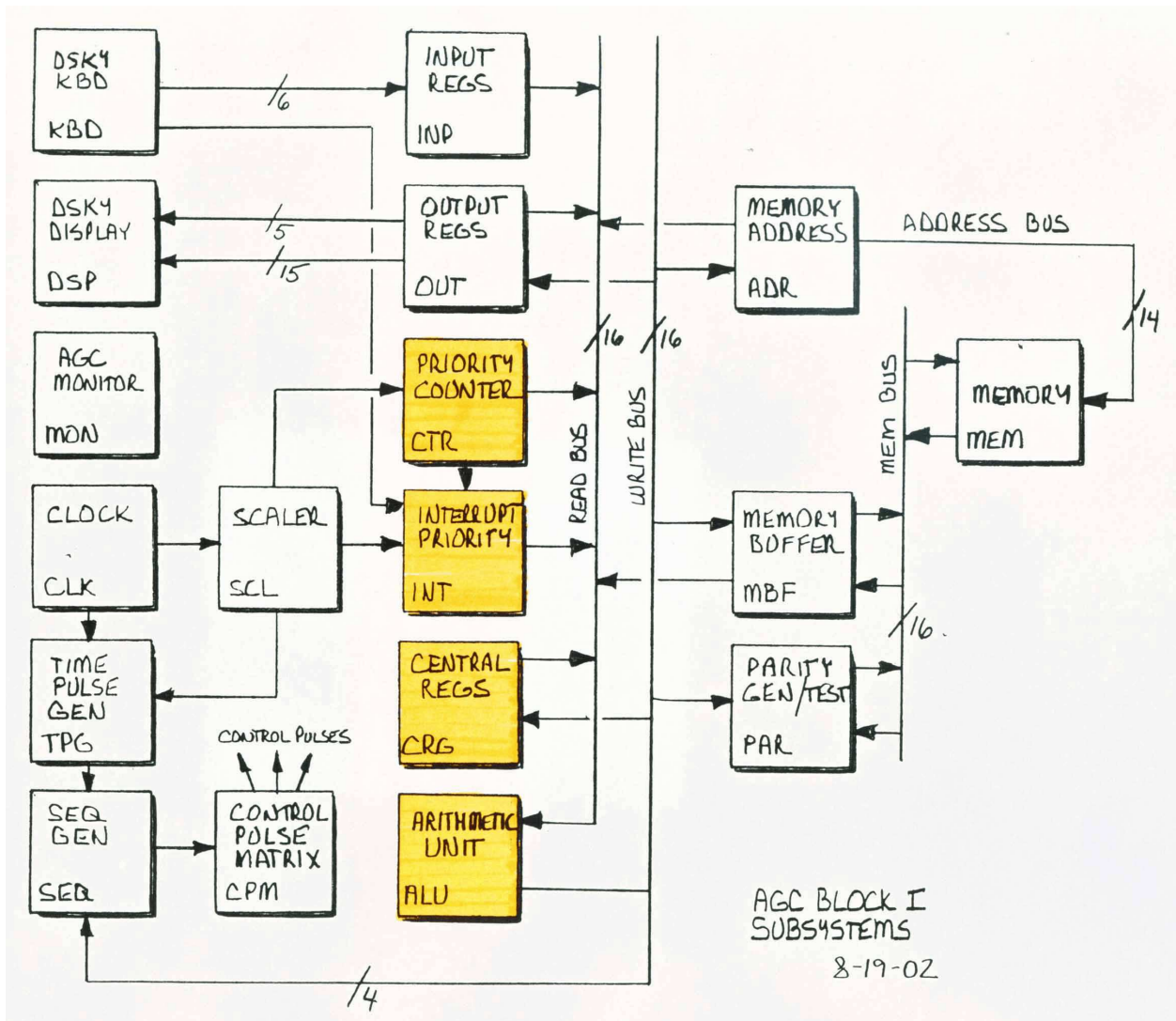
The original AGC had five vectored interrupts. This recreation implements three of them: RUPT1, also called T3RUPT which is used as general-purpose timer by the AGC WAITLIST software; RUPT3, also called T4RUPT or DSRUPT, which is used to update the DSKY display at regular intervals; and RUPT4, also called KERUPT, which is triggered by a key press from the user's keyboard. The AGC responds to each interrupt by temporarily suspending the current program, executing a short interrupt service routine, and then resuming the



interrupted program.

### CTR (Priority Counter)

Twenty memory locations in the original AGC functioned as up/down counters. The counters would increment (PINC) or decrement (MINC) in response to external inputs. Increment or decrement was handled by one 12-step subsequence of microinstructions inserted between any two regular instructions. This replica implements 5 of the counters: OVCTR, an overflow counter incremented or decremented by arithmetic overflow during certain instructions; TIME2 and TIME1, the AGC real-time clock; TIME3, a general purpose timer incremented by a 100Hz signal from the SCALER (SCL); and TIME4, a timer used to update the DSKY display.

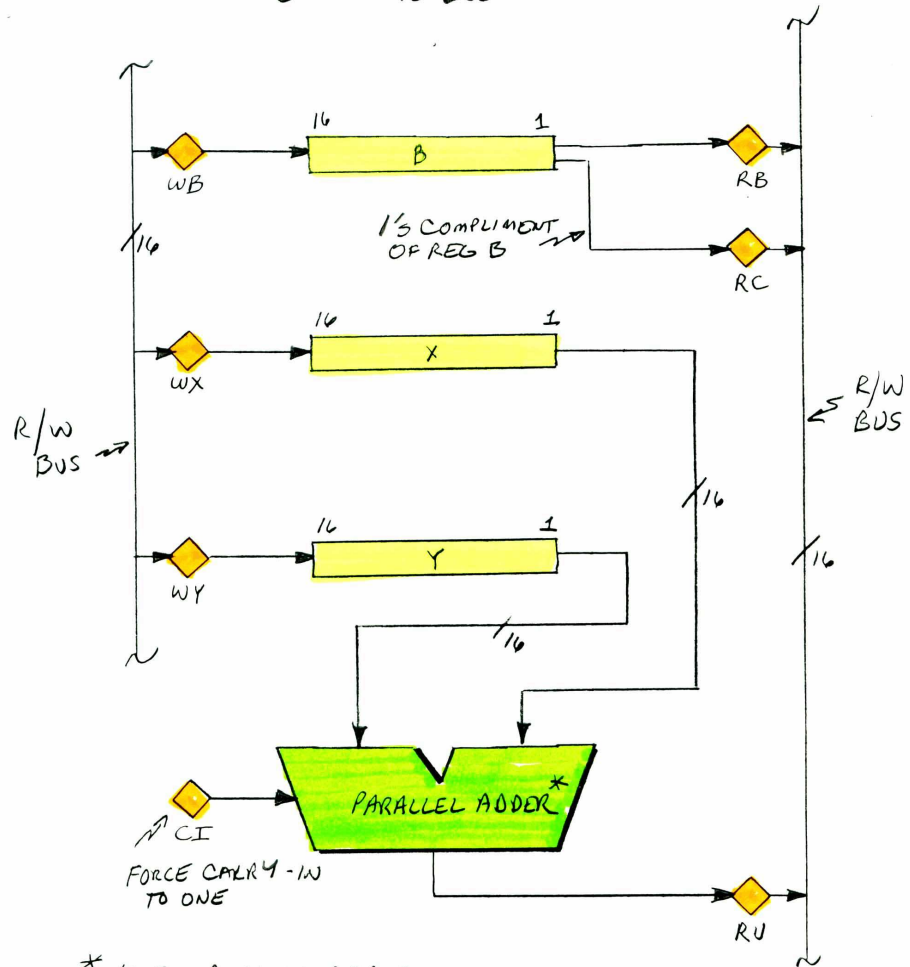


The ALU has a 16-bit parallel adder. The addend and augend are supplied by the X and Y registers. The sum, called the U register although it is not really a register at all, is gated to the read bus through the RU (read U) control signal. The WX and WY control signals copy the contents of the write bus into the X and Y registers.

The B register is loaded from the write bus with the WB (write B) signal. The contents of B are output to the read bus with the RB (read B) signal. The inverted output of B can be gated to the read bus with the RC (read C) signal.

NOTE: ALSO INCLUDES MANY CONTROL PULSES FOR SETTING SUBSETS OF BITS ON THE R/W BUS.

MOST READ PULSES ARE INCLUSIVE OR'ED ON TO THE BUS.

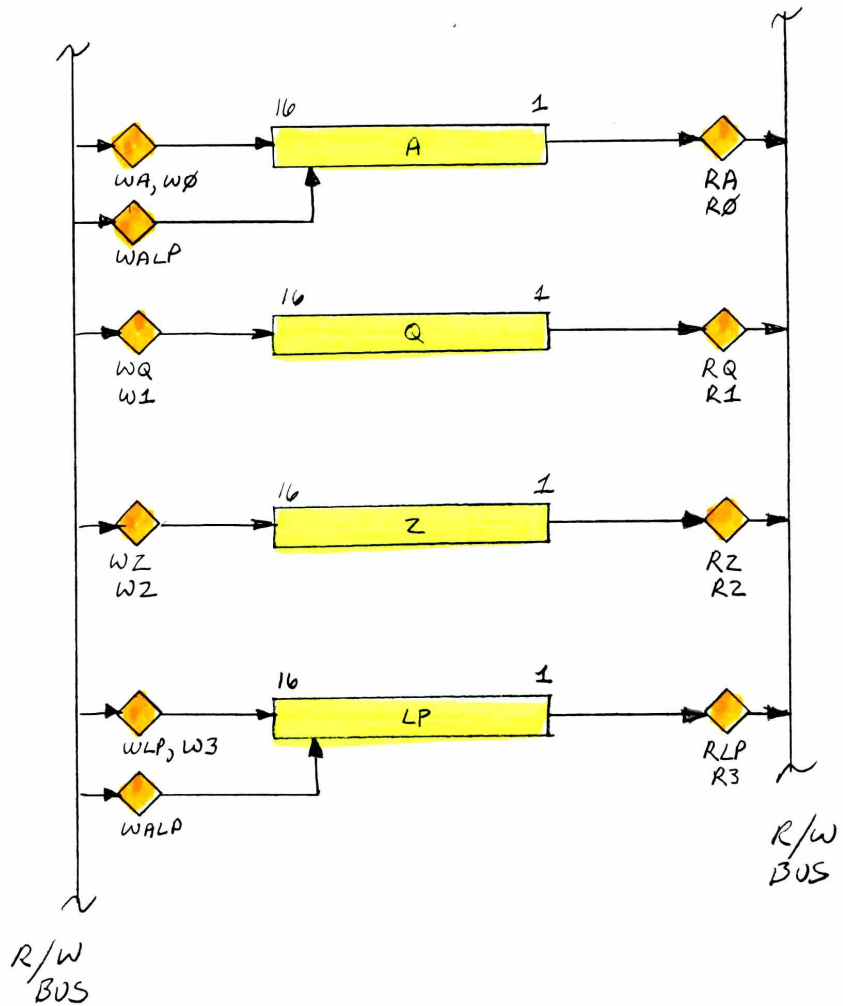


\* 16 BIT PARALLEL ADDER WITH END-AROUND CARRY

The AGC has four user-accessible central registers. The A register is the accumulator; Z is the program counter; Q stores the return address for jumps (TC instruction), and LP stores the lower product (MP instruction).

Central register contents can be output to the read bus by asserting the appropriate read control pulse (RA, RQ, RZ, or RLP). Each register is also mapped to a memory location, with register A mapped to address 0, Q to address 1, Z to 2, and LP to 3. The R0, R1, R2, and R3 control pulses output those registers to the read bus.

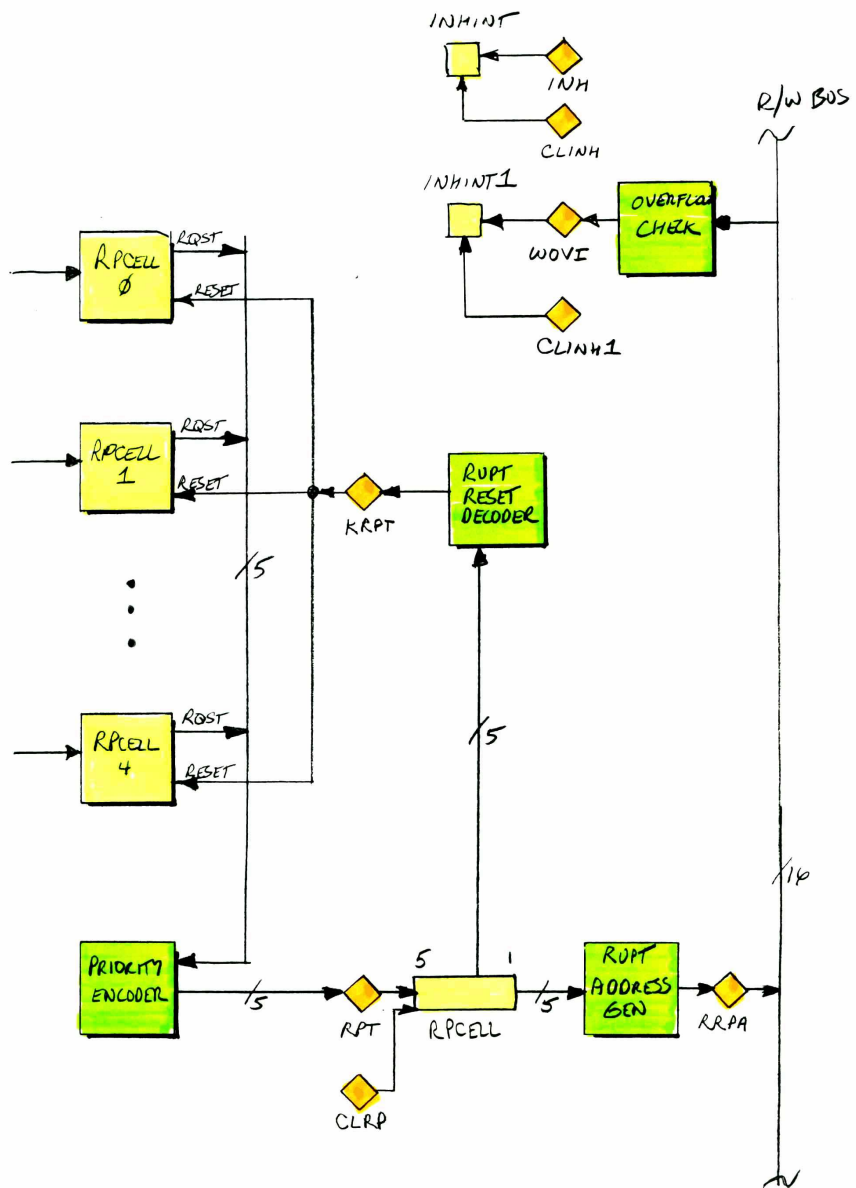
The write bus contents can be loaded into a central register with a write control pulse. WA and WALP load the A register; WQ, the Q register; WZ the Z register; and WLP and WALP the LP register. The W0, W1, W2, and W3 control pulses mapped to memory addresses 0,1,2, and 3 also load those registers.



The interrupt priority subsystem manages vectored interrupts. Five interrupts (0-4) are implemented. Each interrupt is latched by its own "RP cell" flip-flop. Signals from all RP cells feed into a priority encoder; a combinational logic array that outputs the code of the highest priority interrupt in the RP cells. When RPT is asserted, the priority code is latched into RPCELL. This is decoded into an address which is the interrupt vector; the address is written to the read bus when RRPA is asserted.

After the interrupt code has been loaded into RPCELL, asserting KRPT (knock-down RPT) causes the RP cell for that interrupt to be reset. This causes the next highest priority interrupt to be decoded by the priority encoder.

The INHINT and INHINT1 flip-flops inhibit interrupts.

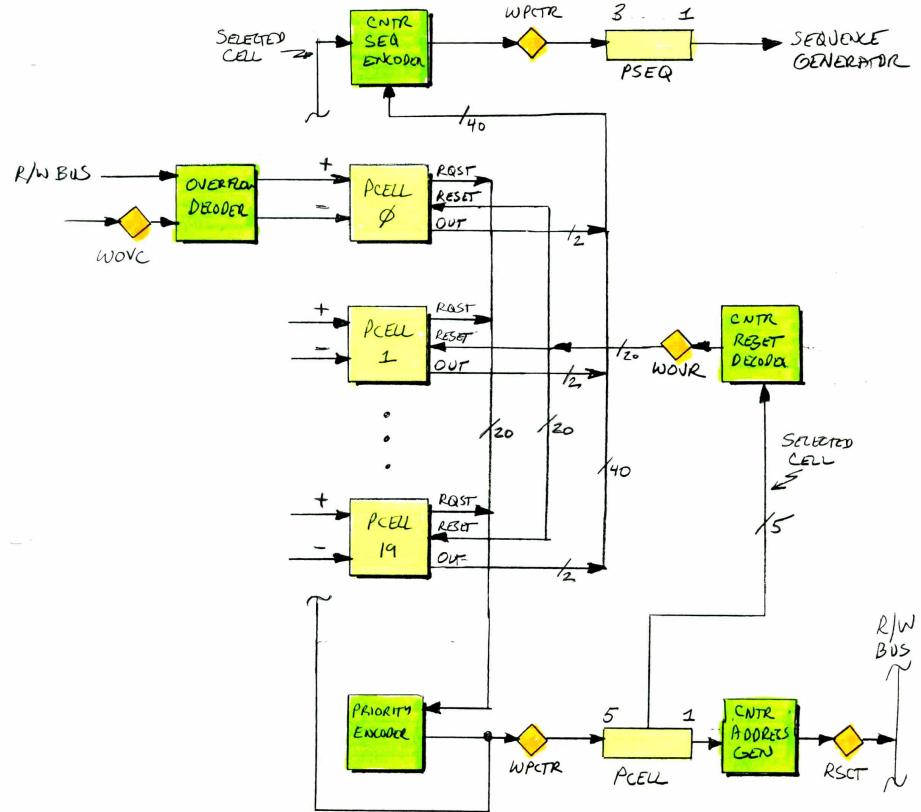


The priority counter logic design is similar to the interrupt subsystem. Up (+) and down (-) count input signals feed into 20 PCELLs, one PCELL for each counter. The PCELLs feed into the priority encoder which outputs the code of the highest priority PCELL having a up or down input set. The PCELL code is written to the PCELL register when WPCTR is asserted.

The PCELL memory address, derived from PCELL, is written to read bus when RSCT is asserted. After the code is latched into the PCELL register, the corresponding PCELL is reset by asserting WOVR.

The up or down code for the selected PCELL is written to PSEQ when WPCTR is asserted. This code feeds to the control logic on the CTL module which selects the PINC (increment) or MINC (decrement) instruction subsequence to bump the priority counter up or down.

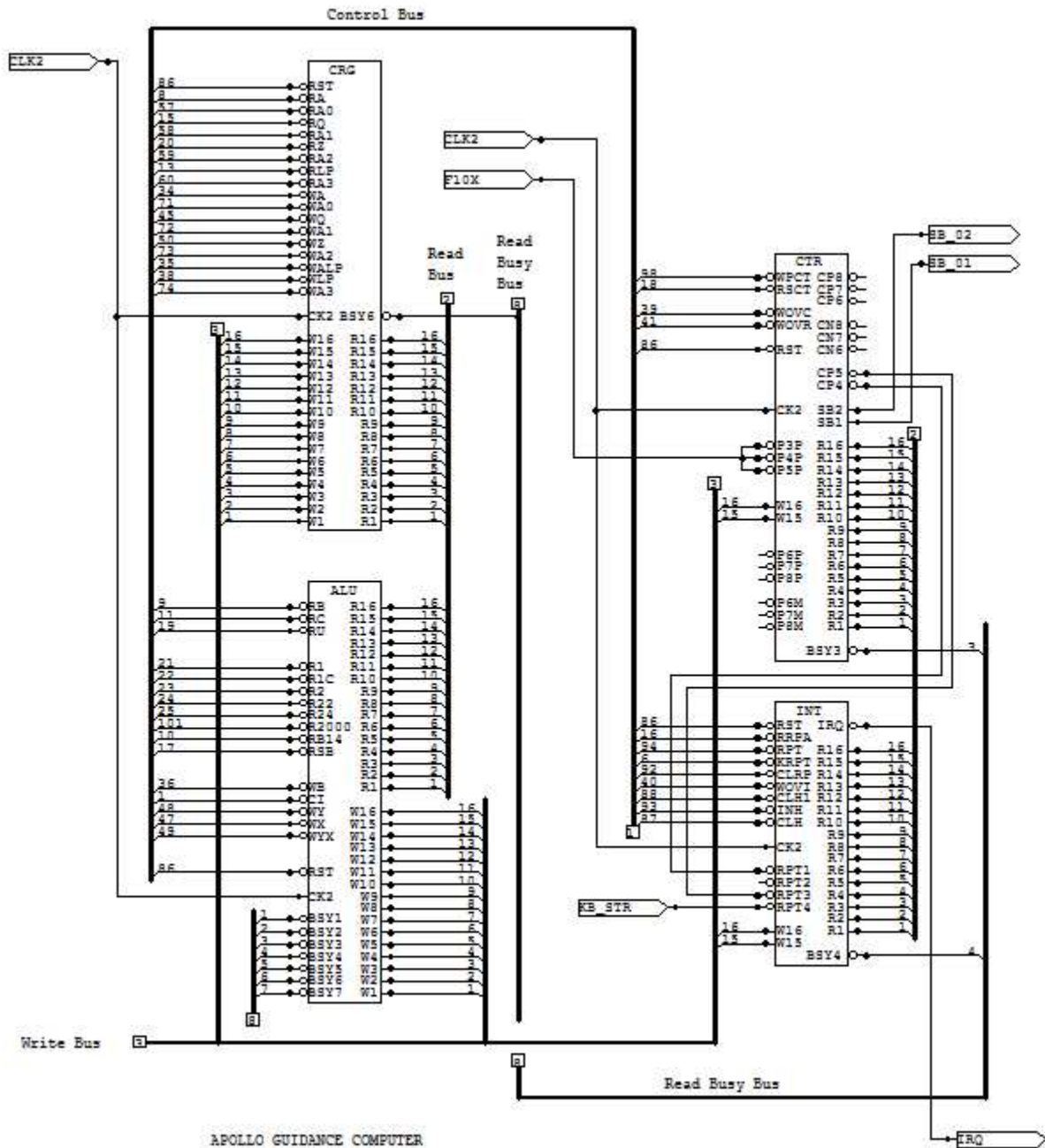
The PINC and MINC subsequences are inserted between normal instruction subsequences. A SHINC subsequence implements a bit-shift which is used to load telemetry bits into the AGC and assemble them into words. SHINC is not implemented in this AGC replica.





# PROC Internal Subsystem Interconnections

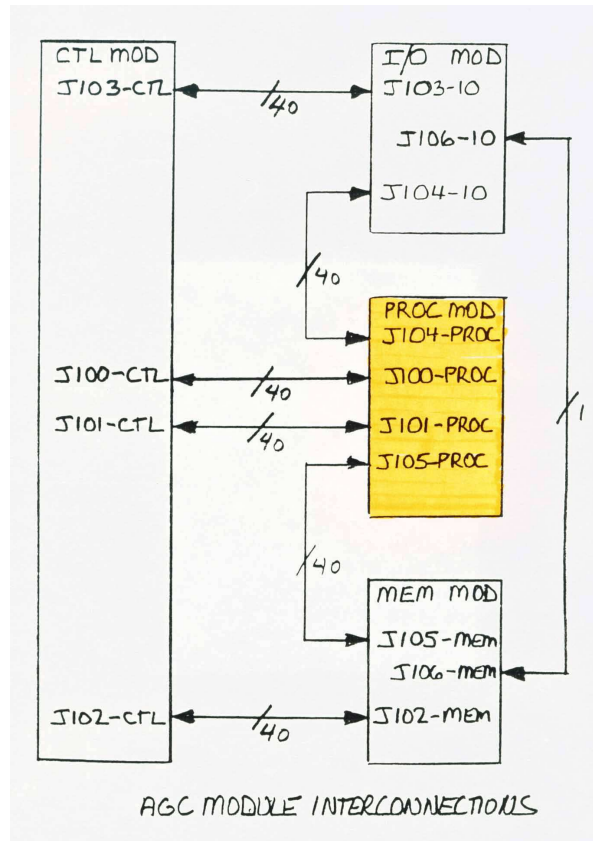
This diagram shows internal interconnections for the subsystems in the PROC module.



APOLLO GUIDANCE COMPUTER  
AGC4 REPLICATED MODULE  
9/2/2003

## PROC Module External Interfaces

The PROC module interfaces to the CTL, MEM, and IO modules through 40-pin IDE ribbon cables.

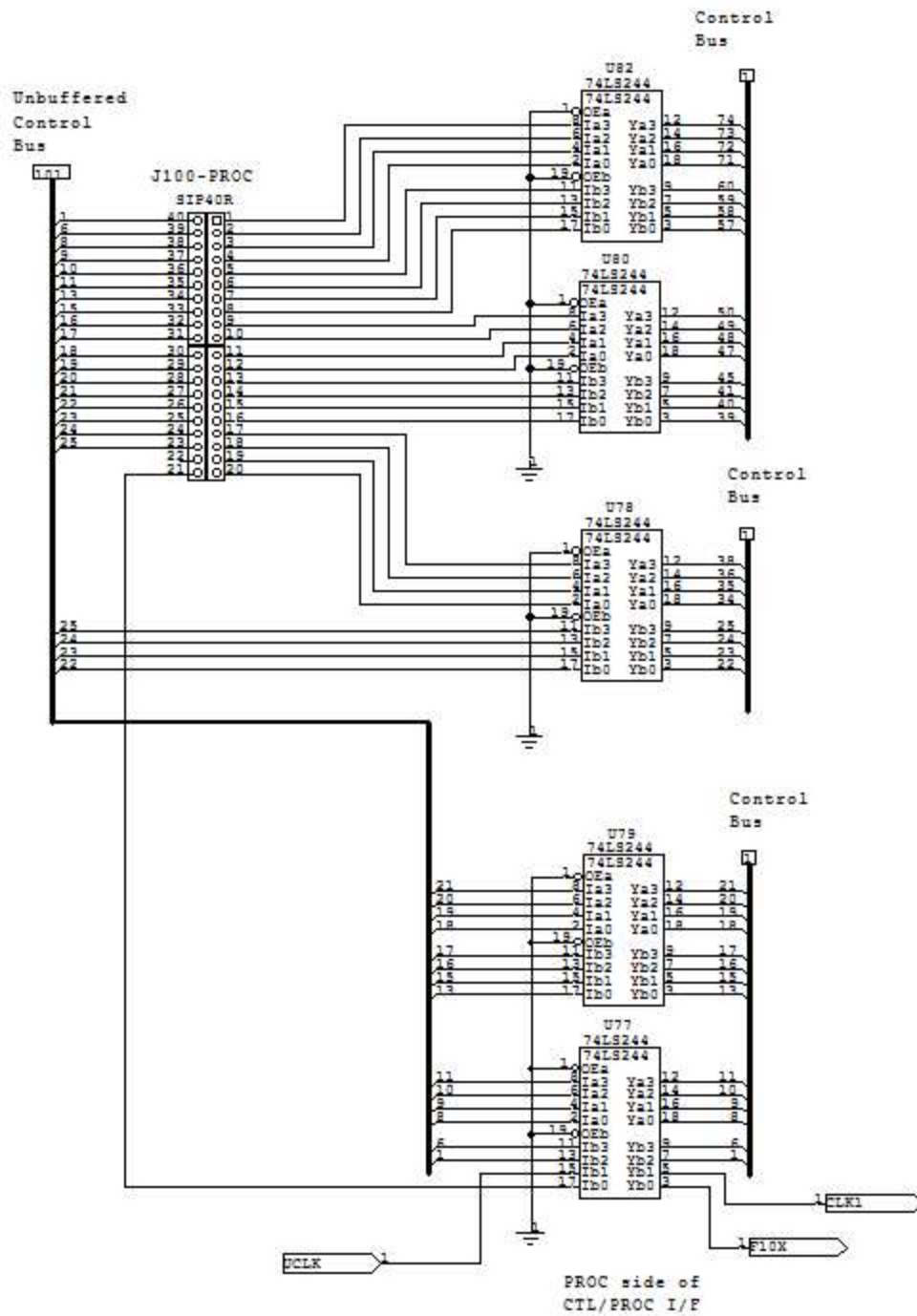


J100-PROC: PROC-to-CTL I/F

J100 is a 40-pin IDE cable that connects the PROC module to the CTL module.

*INPUTS (to PROC):*

<u>PIN</u>	<u>signal</u>	<u>full name</u>	<u>state definition</u>
1	WA3	WRITE ADDR 3 (74)	0=Write reg at address 3 (LP)
2	WA2	WRITE ADDR 2 (73)	0=Write reg at address 2 (Z)
3	WA1	WRITE ADDR 1 (72)	0=Write reg at address 1 (Q)
4	WA0	WRITE ADDR 0 (71)	0=Write reg at address 0 (A)
5	RA3	READ ADDR 3 (60)	0=Read reg at address 3 (LP)
6	RA2	READ ADDR 2 (59)	0=Read reg at address 2 (Z)
7	RA1	READ ADDR 1 (58)	0=Read reg at address 1 (Q)
8	RA0	READ ADDR 0 (57)	0=Read reg at address 0 (A)
9	WZ	WRITE Z (50)	0=Write Z
10	WYx	WRITE Y NO RESET (49)	0=Write Y (do not reset)
11	WY	WRITE Y (48)	0=Write Y
12	WX	WRITE X (47)	0=Write X
13	WQ	WRITE Q (45)	0=Write Q
14	WOVR	WRITE OVF (41)	0=Write overflow
15	WOVI	WRITE OVF RUPT INH (40)	0=Write overflow RUPT inhibit
16	WOVC	WRITE OVF CNTR (39)	0=Write overflow counter
17	WLP	WRITE LP (38)	0=Write LP
18	WB	WRITE B (36)	0=Write B
19	WALP	WRITE A/LP (35)	0=Write A and LP
20	WA	WRITE A (34)	0=Write A
21	F10X	F10 SCALER ONESHOT	1=timed out (100.0 Hz)
23	R24	READ 24 (25)	0=Read 24
24	R22	READ 22 (24)	0=Read 22
25	R2	READ 2 (23)	0=Read 2
26	R1C	READ 1 COMP (22)	0=Read 1 complimented
27	R1	READ 1 (21)	0=Read 1
28	RZ	READ Z (20)	0=Read Z
29	RU	READ U (19)	0=Read sum
30	RSCT	READ CNTR ADDR (18)	0=Read selected counter address
31	RSB	READ SIGN (17)	0=Read sign bit
32	RRPA	READ RUPT ADDR (16)	0=Read RUPT address
33	RQ	READ Q (15)	0=Read Q
34	RLP	READ LP (13)	0=Read LP
35	RC	READ C (11)	0=Read C
36	RB14	READ BIT 14 (10)	0=Read bit 14
37	RB	READ B (9)	0=Read B
38	RA	READ A (8)	0=Read A
39	KRPT	KNOCK DOWN RUPT (6)	0=Knock down Rupt priority
40	CI	SET CARRY IN (1)	0=Carry in



## J101-PROC: PROC-to-CTL I/F

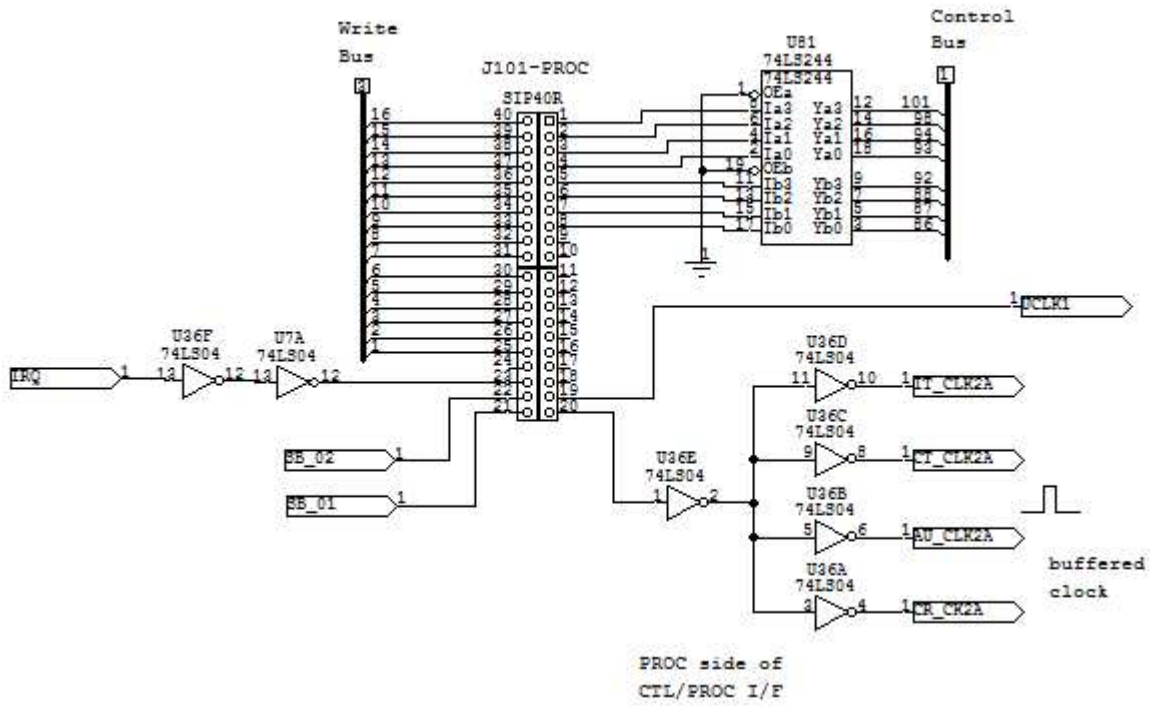
J101 is a 40-pin IDE cable that connects the PROC module to the CTL module.

### *INPUTS (to PROC):*

<u>PIN</u>	<u>signal</u>	<u>full name</u>	<u>state definition</u>
1	R2000	READ 2000 (101)	0=Read 2000
2	WPCTR	WRITE PCTR (98)	0=Write PCTR (latch priority counter seq)
3	RPT	READ RUPT (94)	0=Read RUPT opcode
4	INH	SET INHINT (93)	0=Set INHINT
5	CLRP	CLEAR RPCELL (92)	0=Clear RPCELL
6	CLINH1	CLEAR INHINT1 (88)	0=Clear INHINT1
7	CLINH	CLEAR INHINT (87)	0=Clear INHINT
8	GENRST	GENERAL RESET (86)	0=General Reset
19	CLK1	CLOCK1	1.024 MHz AGC clock 1 (normally low)
20	CLK2	CLOCK2	1.024 MHz AGC clock 2 (normally low)

### *OUTPUTS (from PROC):*

<u>PIN</u>	<u>signal</u>	<u>full name</u>	<u>state definition</u>
21	SB_01	SUB SEL 01	SB_01 is LSB; SB_02 is MSB
22	SB_02	SUB SEL 02	00=no counter; 01=PINC; 10=MINC
23	IRQ	INT RQST	0=interrupt requested.
25	WB_01	WRITE BUS 01	(lsb)
26	WB_02	WRITE BUS 02	
27	WB_03	WRITE BUS 03	
28	WB_04	WRITE BUS 04	
29	WB_05	WRITE BUS 05	
30	WB_06	WRITE BUS 06	
31	WB_07	WRITE BUS 07	
32	WB_08	WRITE BUS 08	
33	WB_09	WRITE BUS 09	
34	WB_10	WRITE BUS 10	
35	WB_11	WRITE BUS 11	
36	WB_12	WRITE BUS 12	
37	WB_13	WRITE BUS 13	
38	WB_14	WRITE BUS 14	
39	WB_15	WRITE BUS 15	US (overflow) bit
40	WB_16	WRITE BUS 16	SG (sign) bit



J104-PROC: PROC-to-IO I/F

J104 is a 40-pin IDE cable that connects the PROC module to the IO module.

*INPUTS (to PROC):*

<u>PIN</u>	<u>signal</u>	<u>full name</u>	<u>state definition</u>
40	RB_01	READ BUS 01	(lsb)
39	RB_02	READ BUS 02	
38	RB_03	READ BUS 03	
37	RB_04	READ BUS 04	
36	RB_05	READ BUS 05	
35	RB_06	READ BUS 06	
34	RB_07	READ BUS 07	
33	RB_08	READ BUS 08	
32	RB_09	READ BUS 09	
31	RB_10	READ BUS 10	
30	RB_11	READ BUS 11	
29	RB_12	READ BUS 12	
28	RB_13	READ BUS 13	
27	RB_14	READ BUS 14	
26	RB_15	READ BUS 15	US (overflow) bit
25	RB_16	READ BUS 16	SG (sign) bit
22	BUSY2	READ BUS BUSY	0=OUT register output to read bus
21	BUSY1	READ BUS BUSY	0=INP register output to read bus
20	KB_STR	KEY STROBE	1=key pressed strobe; to KEYRUPT. Key data is valid on the negative edge of KB_STR. Data is latched until the next keypress.

*OUTPUTS (from PROC):*

<u>PIN</u>	<u>signal</u>	<u>full name</u>	<u>state definition</u>
1	WB_01	WRITE BUS 01	(lsb)
2	WB_02	WRITE BUS 02	
3	WB_03	WRITE BUS 03	
4	WB_04	WRITE BUS 04	
5	WB_05	WRITE BUS 05	
6	WB_06	WRITE BUS 06	
7	WB_07	WRITE BUS 07	
8	WB_08	WRITE BUS 08	
9	WB_09	WRITE BUS 09	
10	WB_10	WRITE BUS 10	
11	WB_11	WRITE BUS 11	
12	WB_12	WRITE BUS 12	
13	WB_13	WRITE BUS 13	
14	WB_14	WRITE BUS 14	
15	WB_15	WRITE BUS 15	US (overflow) bit
16	WB_16	WRITE BUS 16	SG (sign) bit

## J105-PROC: PROC-to-MEM I/F

J105 is a 40-pin IDE cable that connects the PROC module to the MEM module.

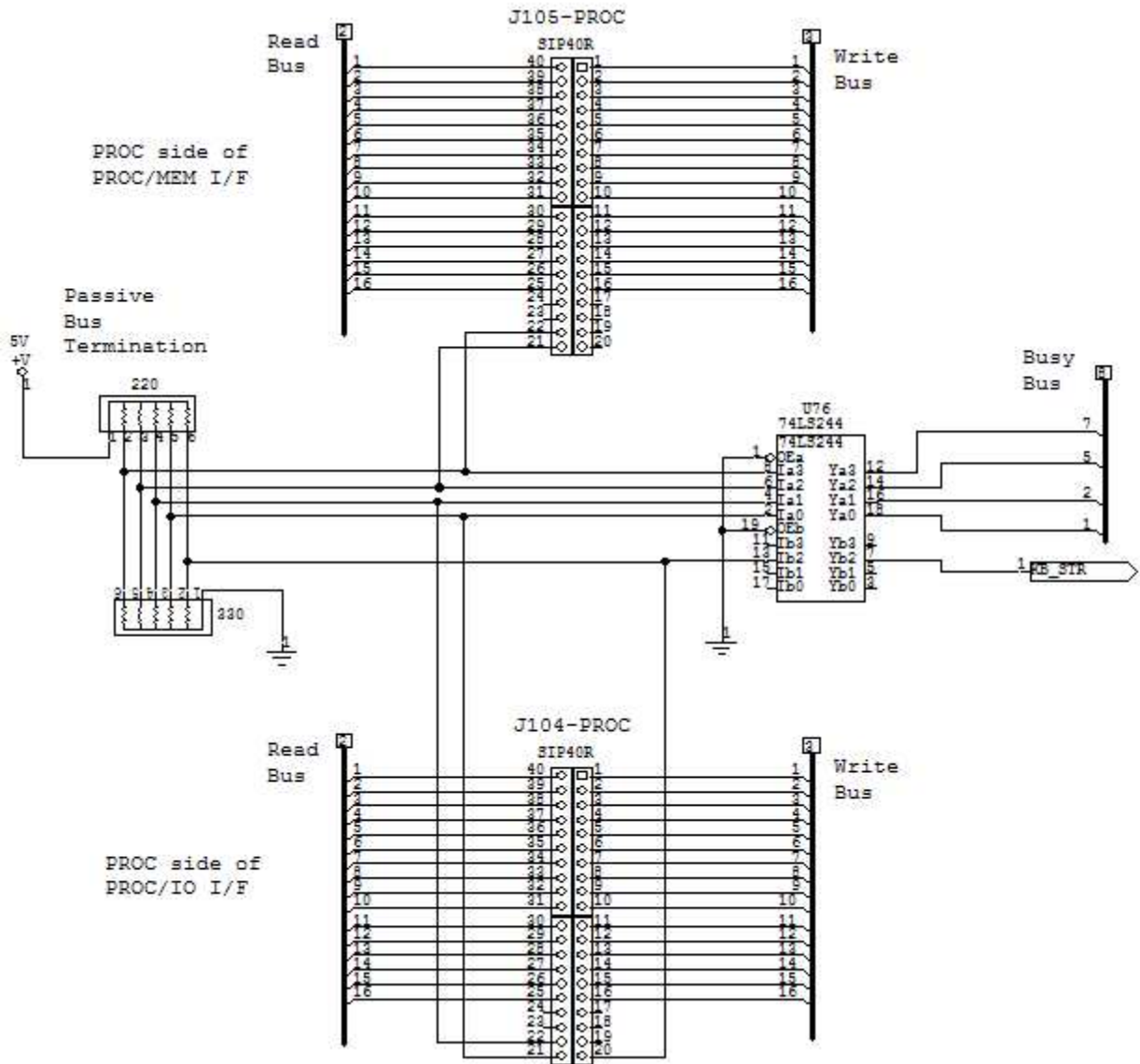
### INPUTS (to PROC):

<u>PIN</u>	<u>signal</u>	<u>full name</u>	<u>state definition</u>
40	RB_01	READ BUS 01	(lsb)
39	RB_02	READ BUS 02	
38	RB_03	READ BUS 03	
37	RB_04	READ BUS 04	
36	RB_05	READ BUS 05	
35	RB_06	READ BUS 06	
34	RB_07	READ BUS 07	
33	RB_08	READ BUS 08	
32	RB_09	READ BUS 09	
31	RB_10	READ BUS 10	
30	RB_11	READ BUS 11	
29	RB_12	READ BUS 12	
28	RB_13	READ BUS 13	
27	RB_14	READ BUS 14	
26	RB_15	READ BUS 15	US (overflow) bit
25	RB_16	READ BUS 16	SG (sign) bit
22	BUSY7	READ BUS BUSY	0=BNK register output enabled to read
bus			
21	BUSY5	READ BUS BUSY	0=G register output enabled to read bus

### OUTPUTS (from PROC):

<u>PIN</u>	<u>signal</u>	<u>full name</u>	<u>state definition</u>
1	WB_01	WRITE BUS 01	(lsb)
2	WB_02	WRITE BUS 02	
3	WB_03	WRITE BUS 03	
4	WB_04	WRITE BUS 04	
5	WB_05	WRITE BUS 05	
6	WB_06	WRITE BUS 06	
7	WB_07	WRITE BUS 07	
8	WB_08	WRITE BUS 08	
9	WB_09	WRITE BUS 09	
10	WB_10	WRITE BUS 10	
11	WB_11	WRITE BUS 11	
12	WB_12	WRITE BUS 12	
13	WB_13	WRITE BUS 13	
14	WB_14	WRITE BUS 14	
15	WB_15	WRITE BUS 15	US (overflow) bit
16	WB_16	WRITE BUS 16	SG (sign) bit





## PROC CONTROL PANEL PUSHBUTTONS



RUPT1	Set the RUPT1 flip-flop (FF). Simulates a TIME3 overflow. Triggers a T3RUPT.
RUPT3	Set the RUPT3 flip-flop (FF). Simulates a TIME4 overflow. Triggers a T4RUPT (DSRUPT).
RUPT4	Set the RUPT4 flip-flop (FF). Simulates a DSKY keypress. Triggers a KEYRUPT.
TIME1	Set the TIME1 flip-flop (FF). Increments the low-order word of the AGC real-time clock.
TIME2	Set the TIME2 flip-flop (FF). Increments the high-order word of the AGC real-time-clock.
TIME3	Set the TIME3 flip-flop (FF). Increments the general purpose timer.
TIME4	Set the TIME4 flip-flop (FF). Increments the display update timer.

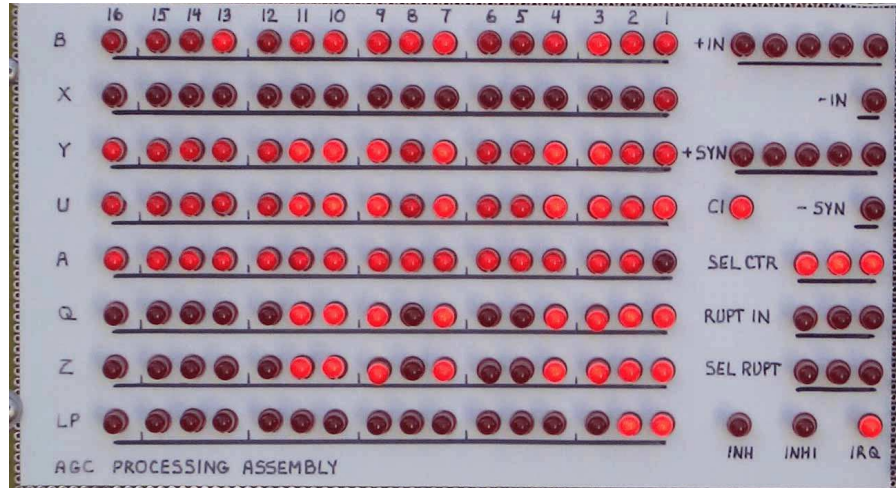
## PROC CONTROL PANEL CONNECTIONS

<u>PIN</u>	<u>signal</u>	<u>state definition</u>
1	RUPT1	GND=set RUPT1 FF
2	RUPT3	GND=set RUPT3 FF
3	RUPT4	GND=set RUPT4 FF
4	TIME1	GND=set TIME1 FF
5	TIME2	GND=set TIME2 FF
6	TIME3	GND=set TIME3 FF
7	TIME4	GND=set TIME4 FF
8	GND	

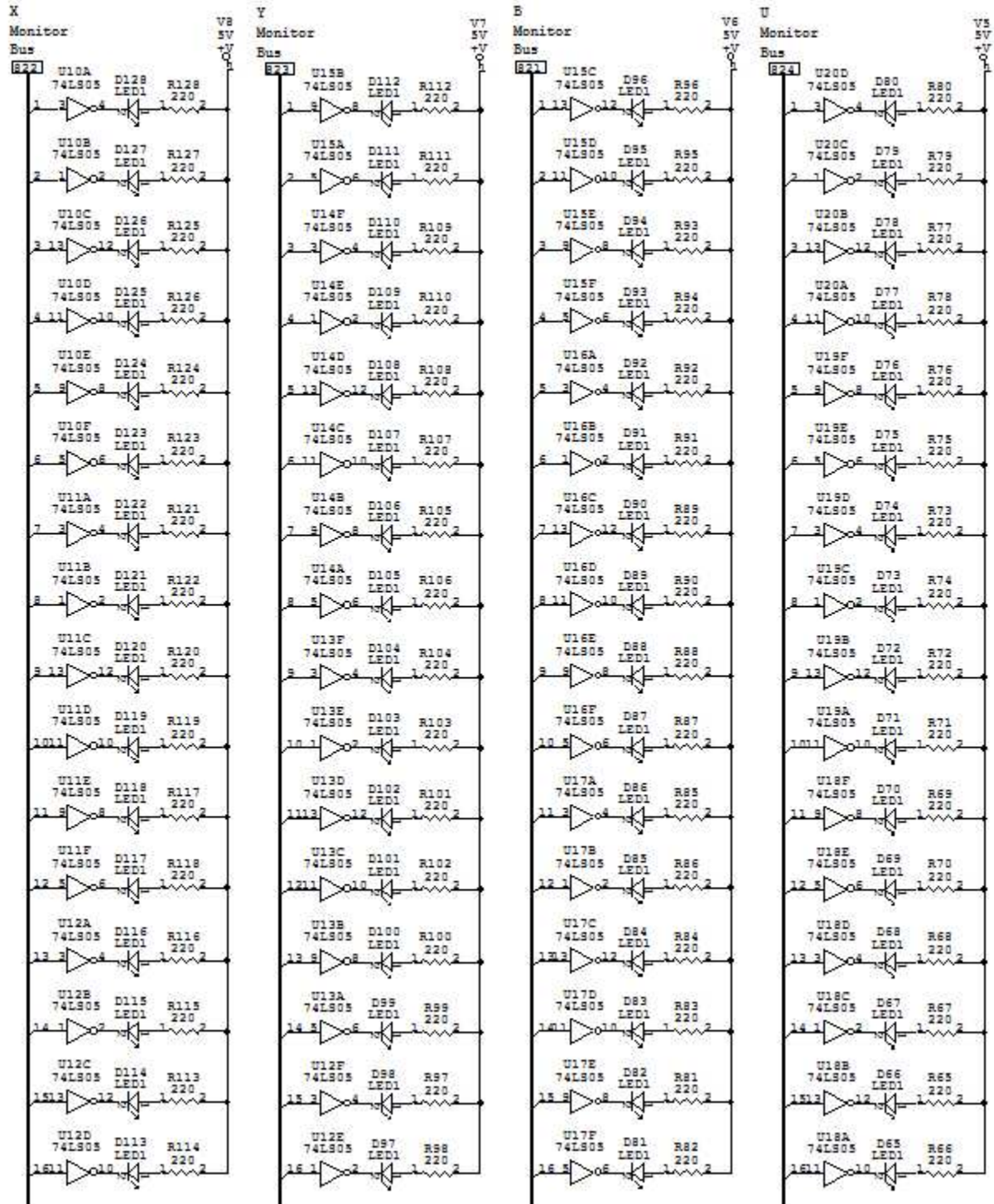
## PROC INDICATORS

The PROC module has a panel of indicator lamps (LEDs) to show the state of PROC registers and critical logic signals.

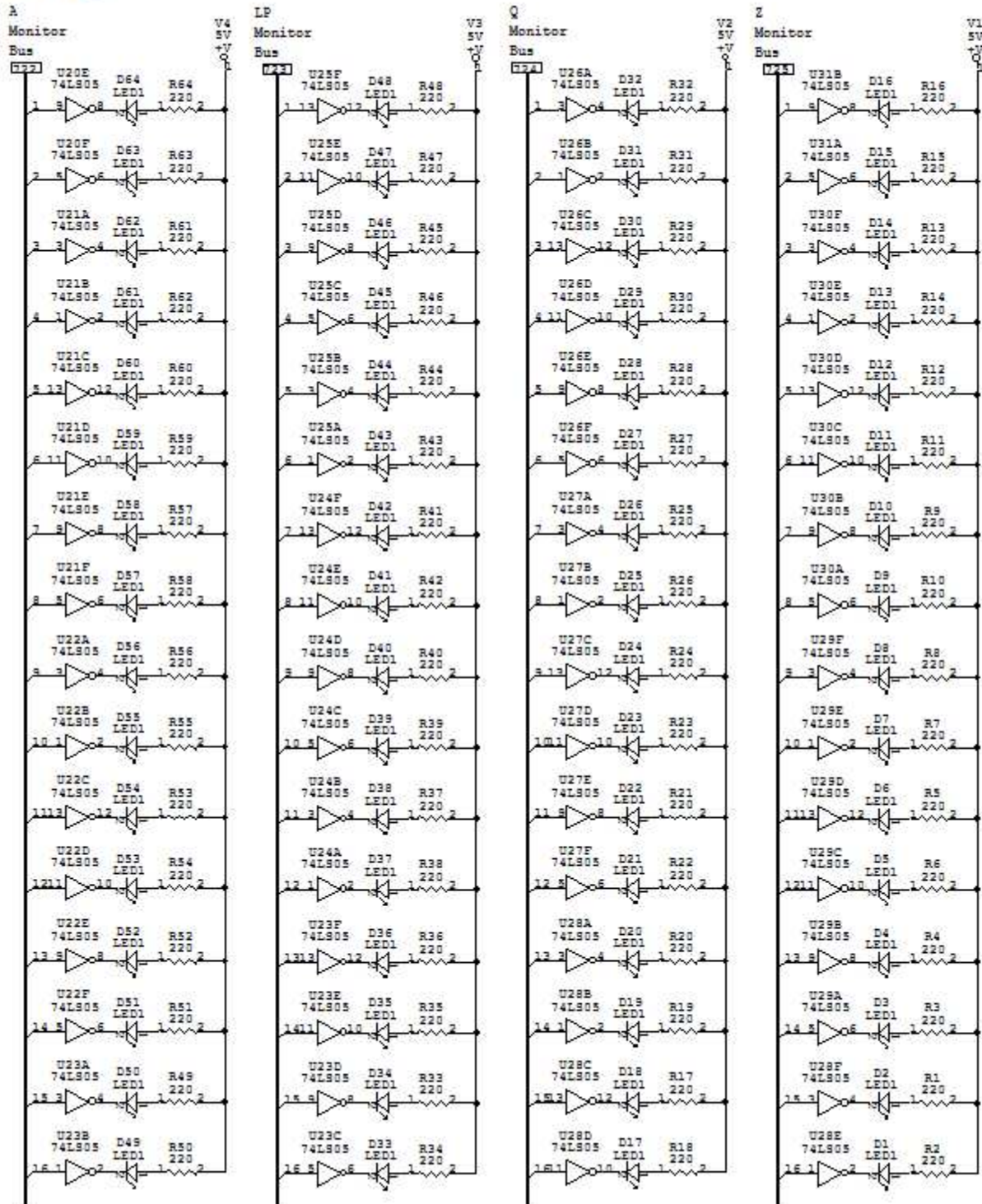
These indicator lamps show the current state of all registers and some additional, important logic signals produced by the PROC module. AGC numbers are represented in octal, so all register lamps are in groups of three. At the time the photo was taken the AGC was running the COLOSSUS 249 flight software load, executing Verb 16, Noun 36: a monitor verb which displays the AGC real time clock.

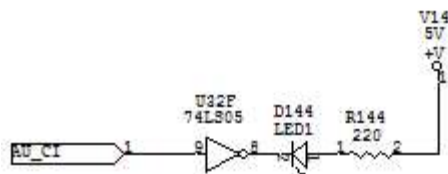
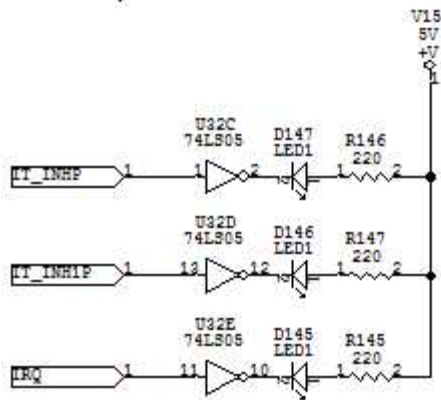
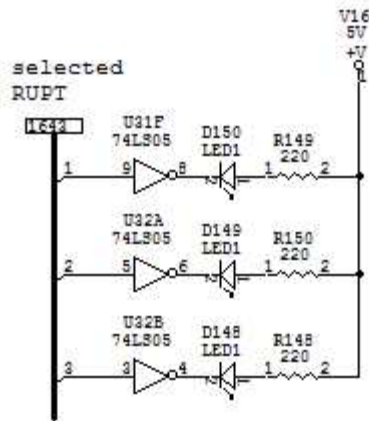
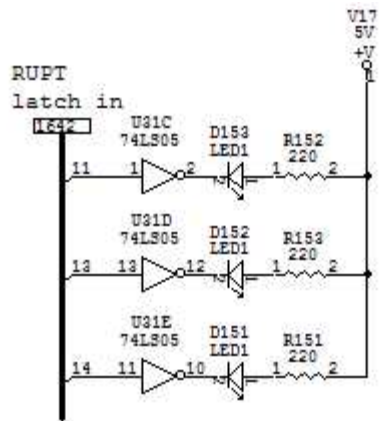


PROC ASSEMBLY  
INDICATORS #1

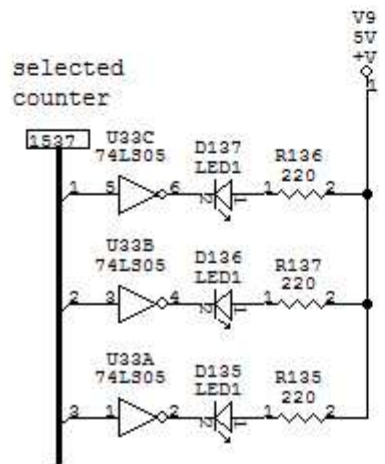
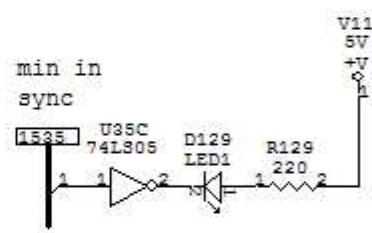
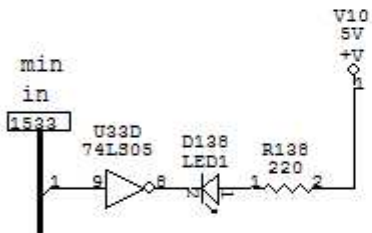
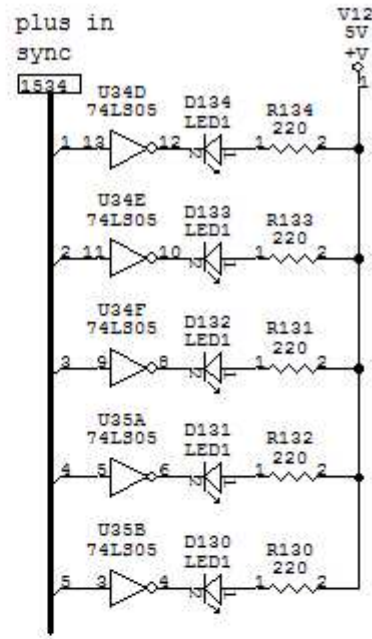
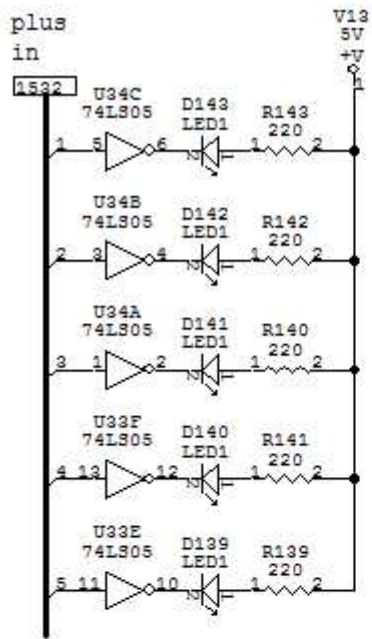


PROC ASSEMBLY  
INDICATORS #2





PROC ASSEMBLY  
INDICATORS #3

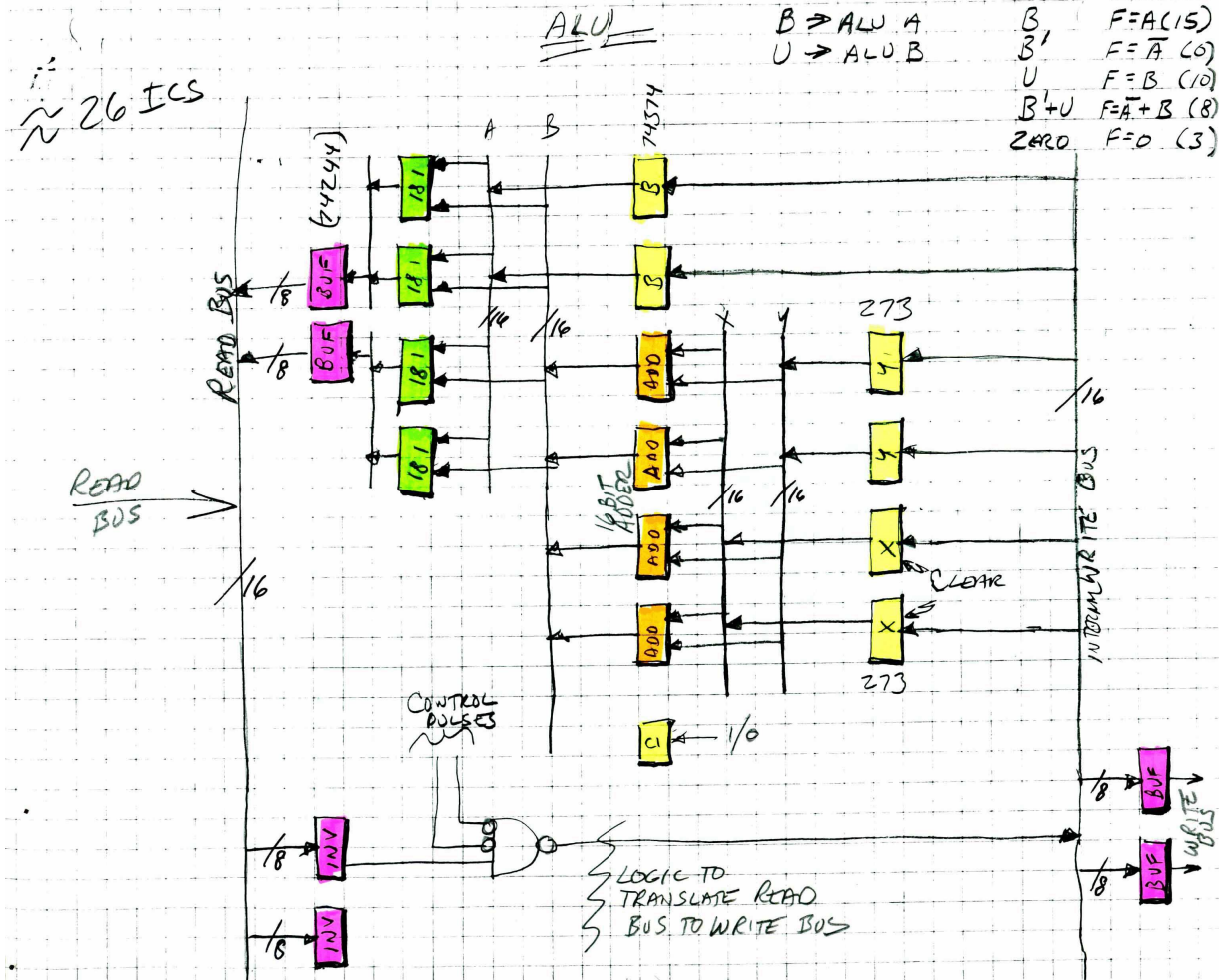


PROC ASSEMBLY  
INDICATORS #4



# ALU (Arithmetic Logic Unit)

My earliest architectural representation of the ALU logic is shown below:



The ALU contains the 16-bit ADDER (colored orange in the diagram) which performs 1's complement arithmetic, and increments the program counter (Z register). Each orange box is a 4-bit parallel adder; collectively, they add 16 bits. The ADDER uses the X, Y, and U registers:

- X: the 16-bit extension register (2 8-bit registers in yellow) that holds one of two inputs to the ADDER.
- Y: the 16-bit extension register (also in yellow) that holds the other input to the ADDER.
- U: the ADDER output (the 1's complement sum of the contents of registers X and Y). Outputs to the bus labeled "B" on the diagram.

The ALU also contains the B and C registers:

- B: a general-purpose buffer register (shown as 2 8-bit registers in yellow), also used to pre-fetch the next instruction. At the start of the next instruction sequence, the upper bits of B (containing the next op code) are copied to the SQ register (in CTL), and the lower bits (the address) are copied to the S register in (MEM). Output to the bus labeled "A" on the diagram.
- C: not a separate register, but the 1's complement of B.

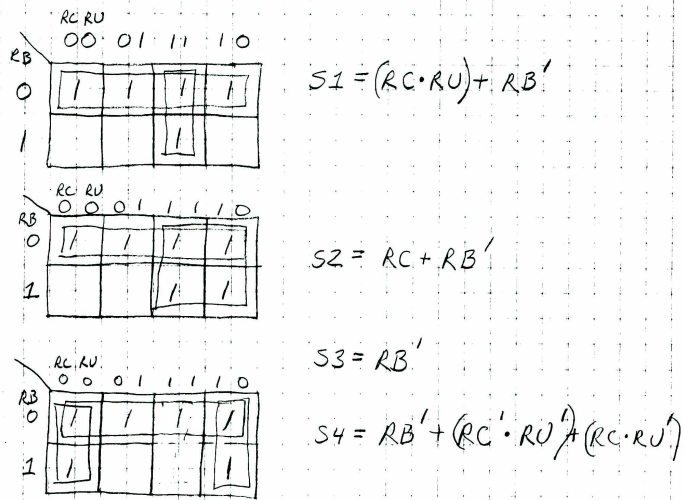
The ALU contains logic (using 74LS181 ALU chips, shown in green) to do any of the following: select the B register; select the complement of the B register (the "C" register); select the U register; select the C register OR'ed with U, or select logical zero. Those logic functions, needed for AGC operation, are shown in the upper right corner of the diagram. The outputs of the 74LS181 selector are gated through a buffer to the read bus.

ALU - 74181 DATA SELECTOR LOGIC

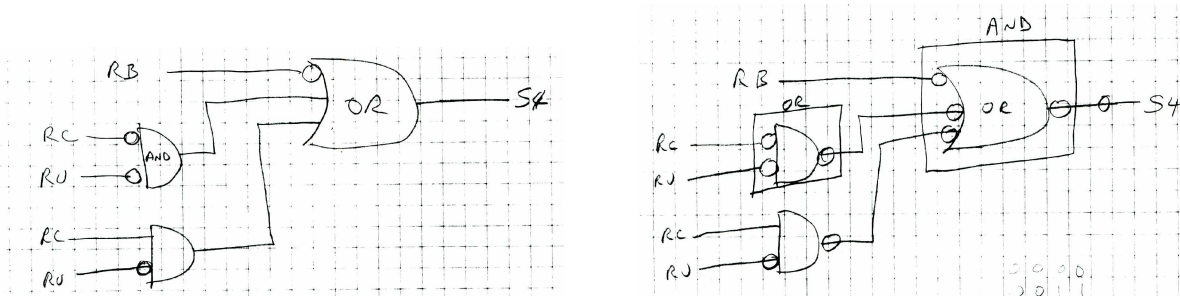
REG B → ALU.A } SELECT INPUT ASSIGNMENTS  
REG U → ALU.B }

INPUTS (ACTIVE LOW)			S				M=H	REGISTER SELECTED
RB	RC	RU	4	3	2	1	(LOGIC EQUATIONS)	
0	0	0	1	1	1	1	F=A (15)	B
0	0	1	1	1	1	1	F=A (15)	B
0	1	0	1	1	1	1	F=A (15)	B
0	1	1	1	1	1	1	F=A (15)	B
1	0	0	1	0	0	0	F=A'+B (8)	B'+U
1	0	1	0	0	0	0	F=A'(0)	B'
1	1	0	1	0	1	0	F=B (10)	U
1	1	1	0	0	1	1	F=0 (3)	ZERO

The original AGC could inclusive OR the outputs of any combination of registers onto the bus, but the control module only used this feature for the B/C and U registers. One of the uses involved the MASK instruction, which is a logical AND: DeMorgan's theorem was used to implement the equivalent of a logical AND by inverting operands through the C register, performing a logical OR through the bus, and then inverting the result.



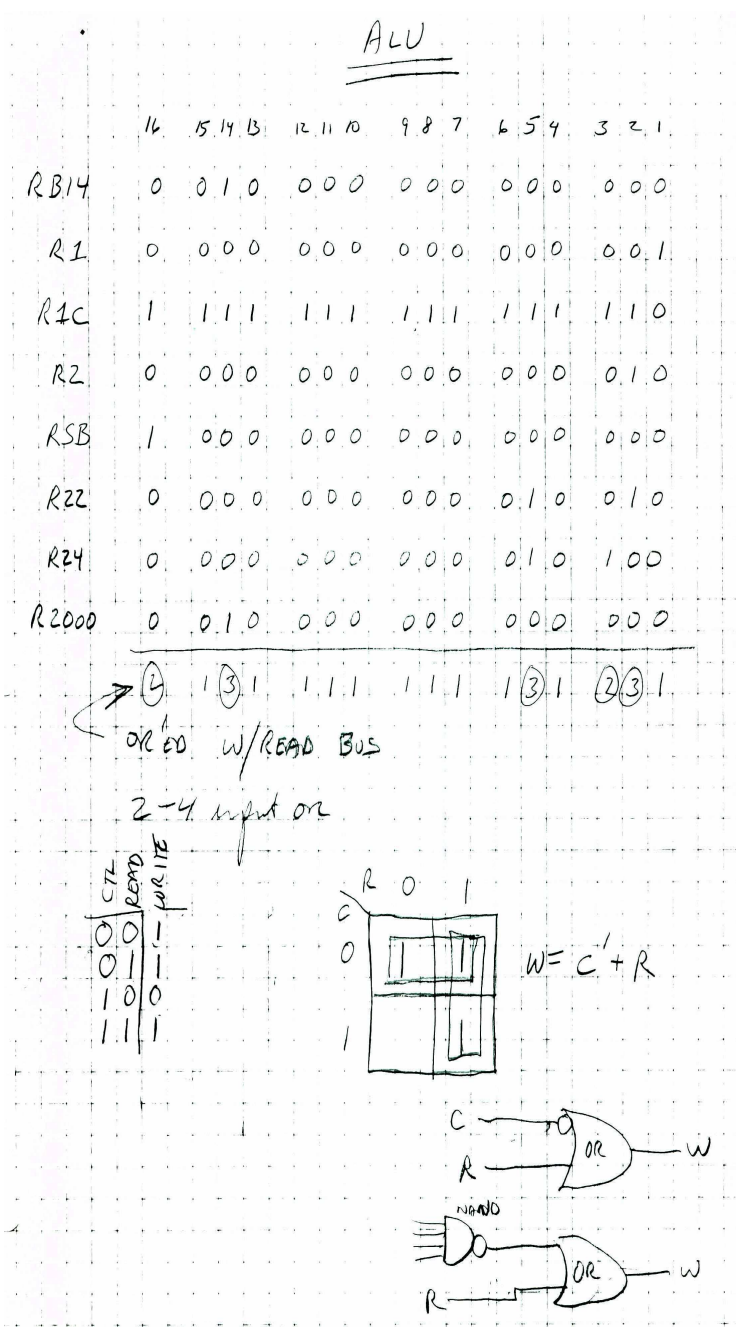
The 74LS181 logic functions are gated to the read bus through the data selector logic, shown in the adjacent diagram.



The ALU also transfers data from the READ bus to the WRITE bus for register-to-register data moves: Data is output from the source register onto the READ bus, then transferred from the READ bus to the WRITE bus through the ALU, and finally loaded from the WRITE bus into the destination register.

The logic that translates the READ bus to the WRITE bus can also force the data lines of the WRITE bus to specific states to gate various arithmetic constants onto the bus. The accompanying diagram shows the control signal on the left (RB14, R1, etc) and to the right is the bit pattern that's OR'ed onto the bus when the signal is asserted. The default state of the read bus is logical zero, so if no other read signal is asserted, the number that appears on the bus is the constant; otherwise, it's the constant inclusive OR'ed with the contents of the READ bus.

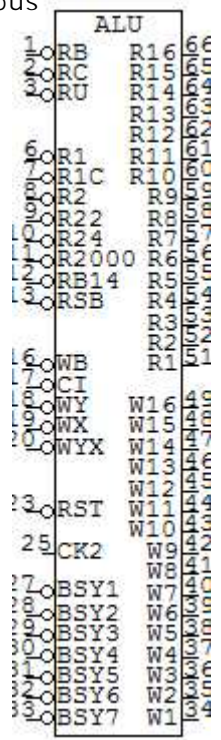
The ALU contains READ bus control logic. Registers in the MEM module, the central registers in the PROC module, and the ALU all interface to the READ bus through tri-state buffers. These buffers are normally in the high-impedance state, but the control module (CTL) issues READ control signals to output the contents of specific registers to the READ bus at certain times. Only one register should be gated onto the READ bus at any given time.



When no READ control signals are asserted, the ALU gates its output onto the READ bus by default. When the control module gates a register onto the READ bus, a BUSY signal is sent to the ALU module, which causes the ALU to inhibit its output. Because of propagation delays, the ALU might continue to output to the READ bus for a brief time while another register is also gated to the bus. To prevent this, all output to the read bus is inhibited during CLK1. This gives the control signals, which transition on the leading edge of CLK1, enough setup time to resolve the conflict.

## ALU INPUTS:

<u>I/F</u>	<u>signal</u>	<u>full name</u>	<u>state definition</u>
CLK:	CLK1	CLOCK 1	1=read bus setup; inhibit read bus out data transfer occurs on falling edge
	CLK2	CLOCK 2	
CPM:	RB	READ B	0=output B register to write bus
	RC	READ C	0=output comp of reg B (C) to write bus
	RU	READ G	0=output U register to write bus
	R1	READ OCTAL 1	0=incl OR 000001 w/write bus
	R1C	READ OCTAL -1	0=incl OR 177776 w/write bus
	R2	READ OCTAL 2	0=incl OR 000002 w/write bus
	R22	READ OCTAL 22	0=incl OR 000022 w/write bus
	R24	READ OCTAL 24	0=incl OR 000024 w/write bus
	R2000	READ OCTAL 2000	0=incl OR 002000 w/write bus
	RB14	READ BIT14	0=incl OR 020000 w/write bus
	RSB	READ SIGN BIT	0=incl OR 100000 w/write bus
	WB	WRITE B	0=write into B from write bus
	CI	WRITE CI	0=set carry register to 1
	WY	WRITE Y	0=write Y
	WX	WRITE X	0=write into X from write bus
	WYX	WRITE Y	0=write into Y from write bus
RBUS:	RB_01	READ BUS 01	US (overflow) bit for read bus
	...		
	RB_14	READ BUS 14	
	RB_15	READ BUS 15	SG (sign) bit for read bus
	RB_16	READ BUS 16	
INP:	BUSY1	READ BUS BUSY	0=valid data from INP on read bus
OUT:	BUSY2	READ BUS BUSY	0=valid data from OUT on read bus
CTR:	BUSY3	READ BUS BUSY	0=valid data from CTR on read bus
INT:	BUSY4	READ BUS BUSY	0=valid data from INT on read bus
MBF:	BUSY5	READ BUS BUSY	0=valid data from MBF on read bus
CRG:	BUSY6	READ BUS BUSY	0=valid data from CRG on read bus
ADR:	BUSY7	READ BUS BUSY	0=valid data from ADR on read bus



## MBF OUTPUTS:

<u>I/F</u>	<u>signal</u>	<u>full name</u>	<u>state definition</u>
------------	---------------	------------------	-------------------------

WBUS:

WB\_01      WRITE BUS 01

...

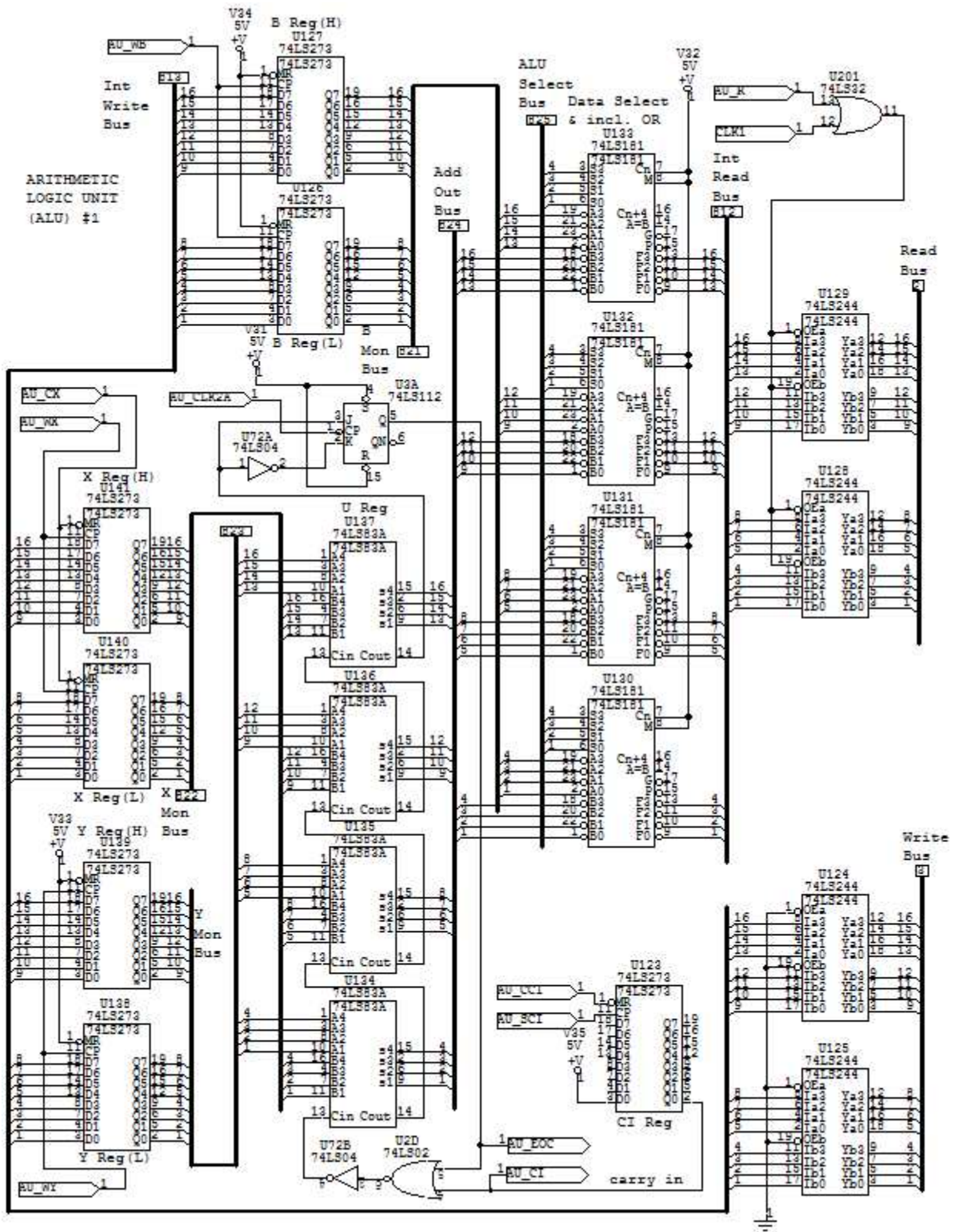
WB\_14      WRITE BUS 14

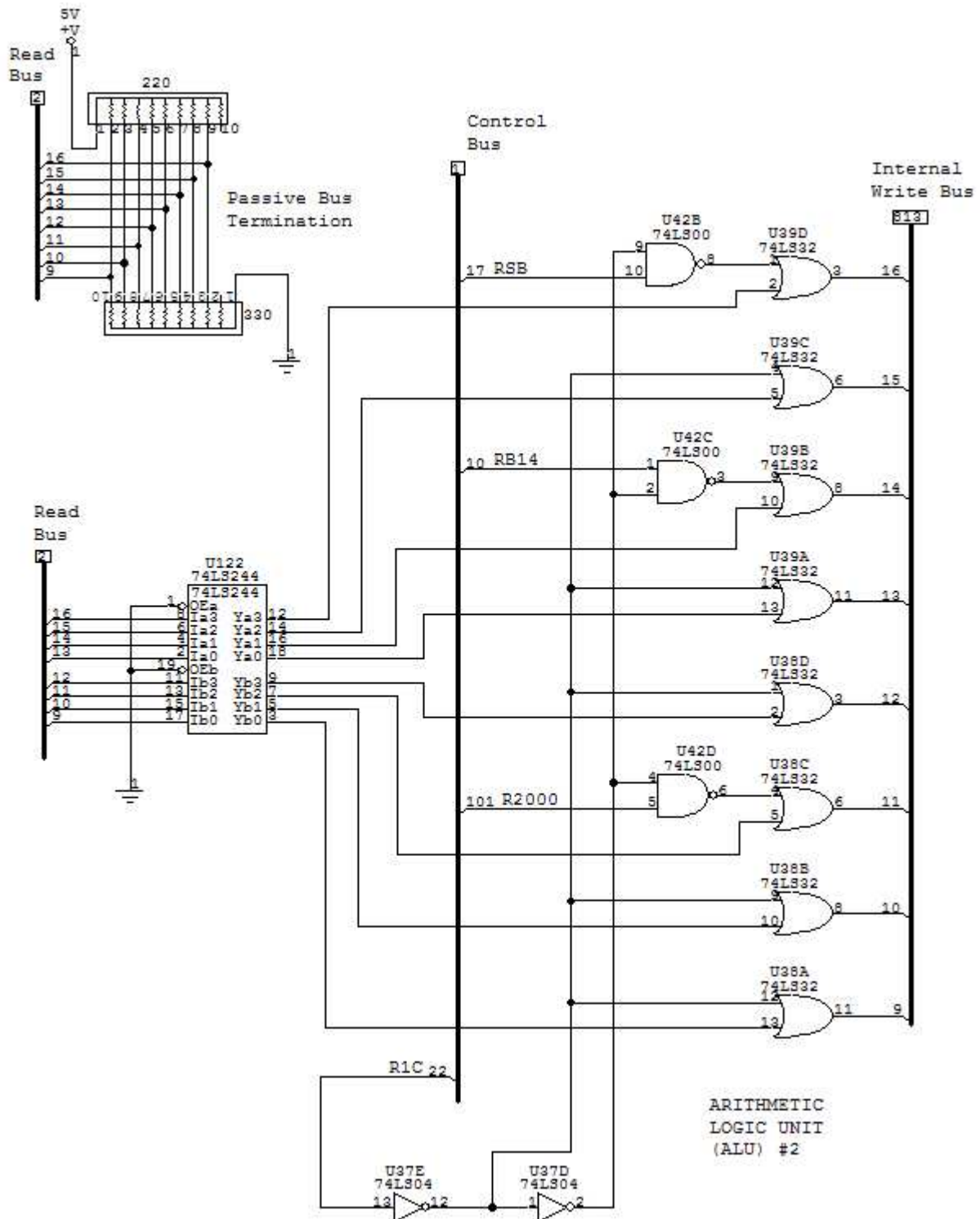
WB\_15      WRITE BUS 15

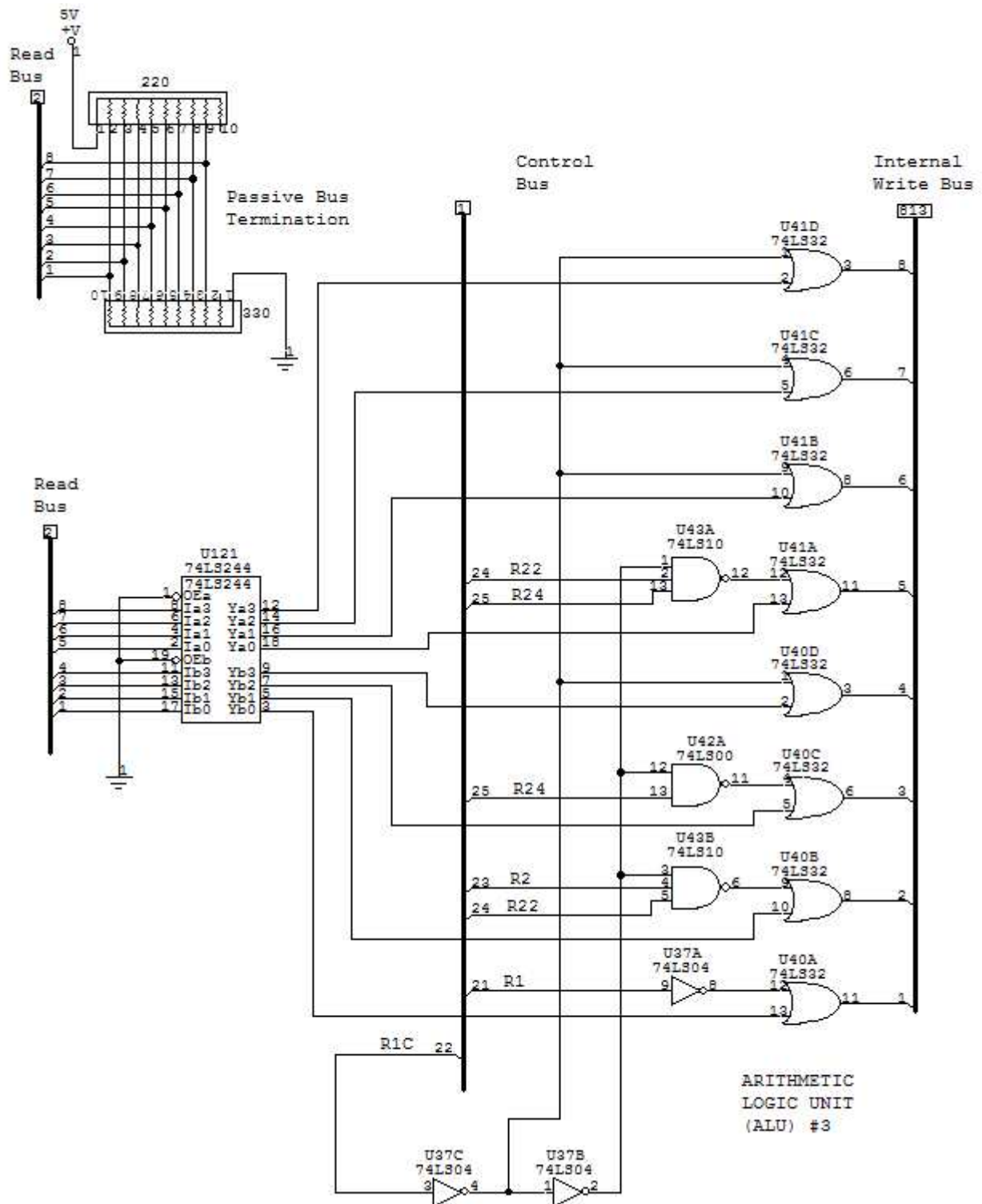
WB\_16      WRITE BUS 16

US (overflow) bit for write bus

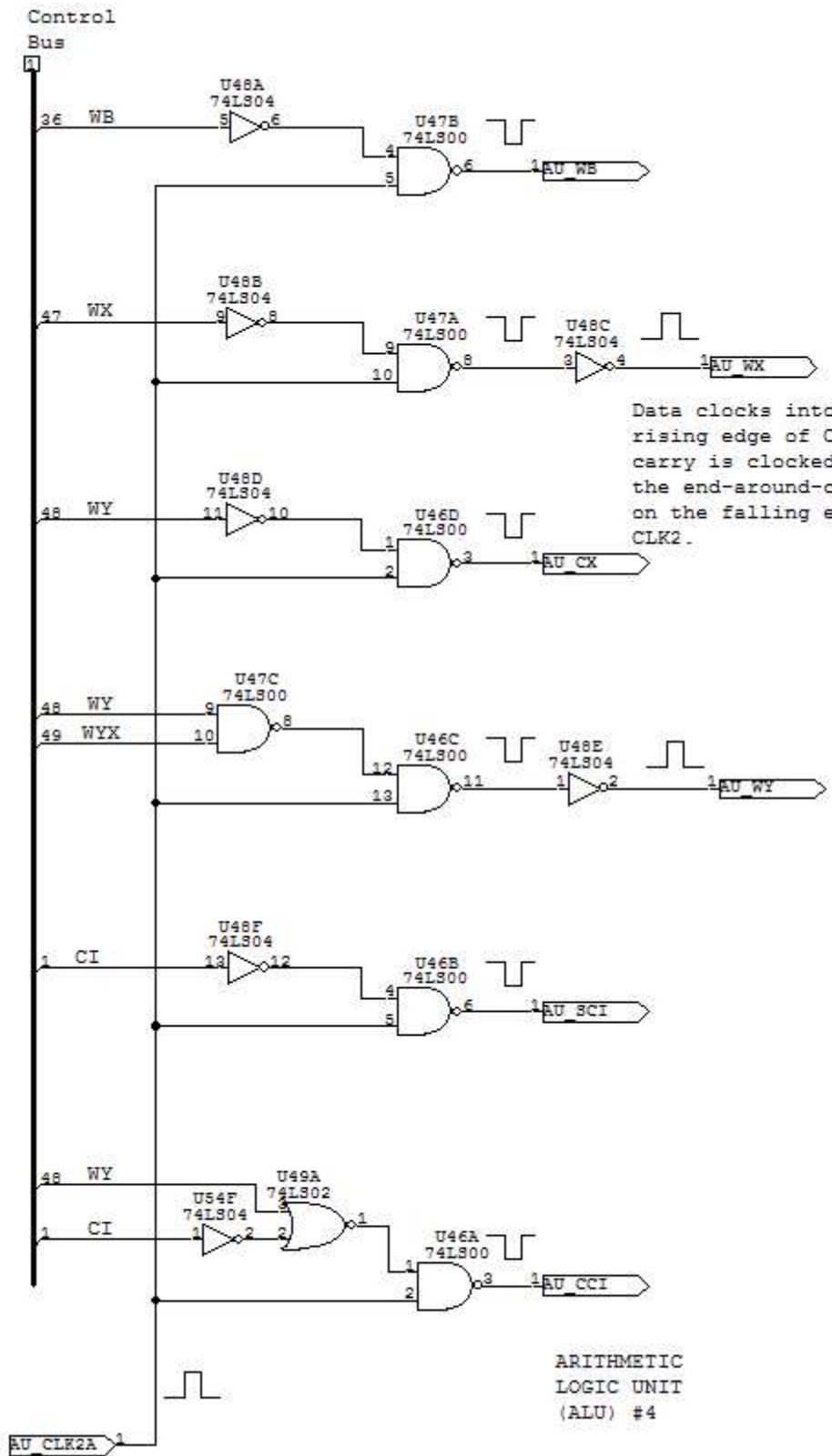
SG (sign) bit for write bus

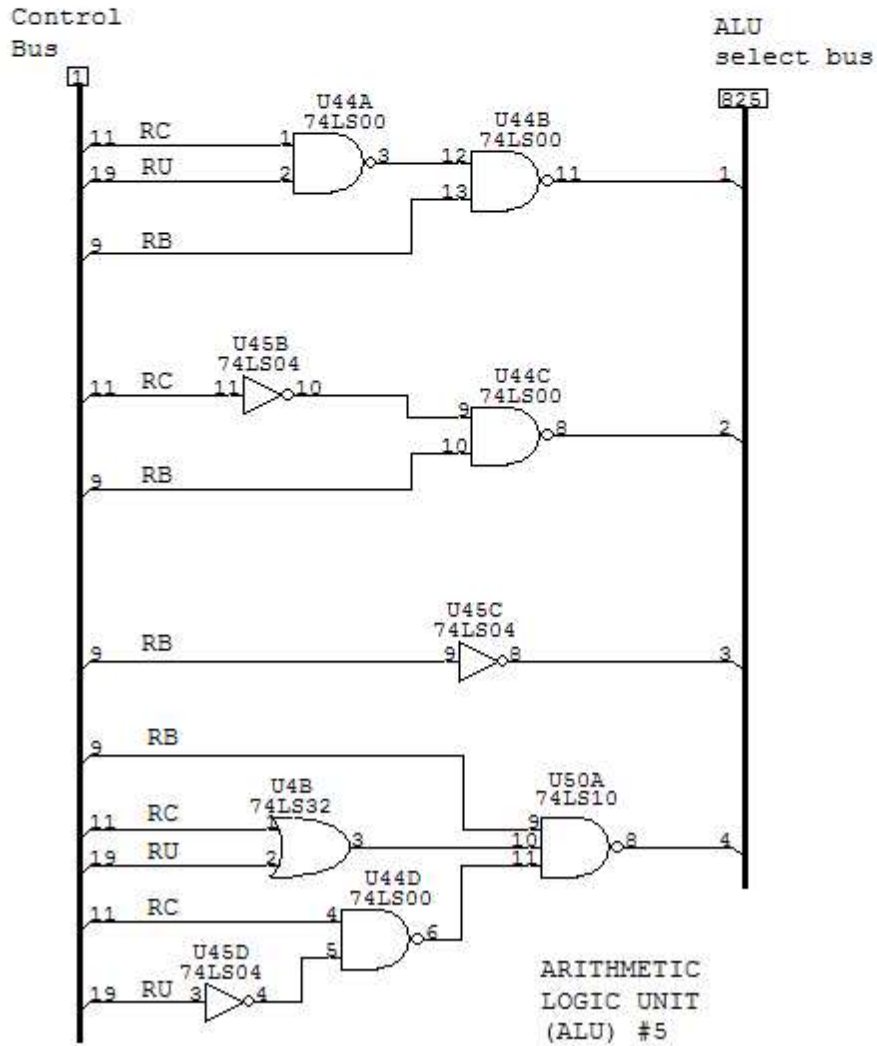
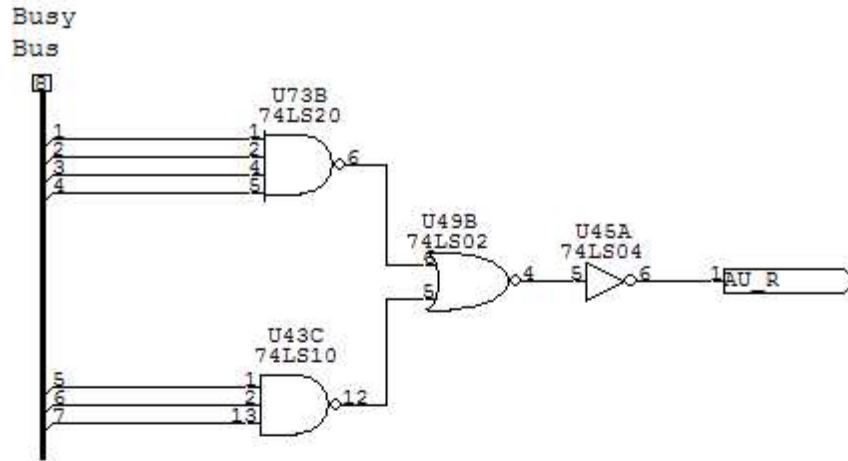








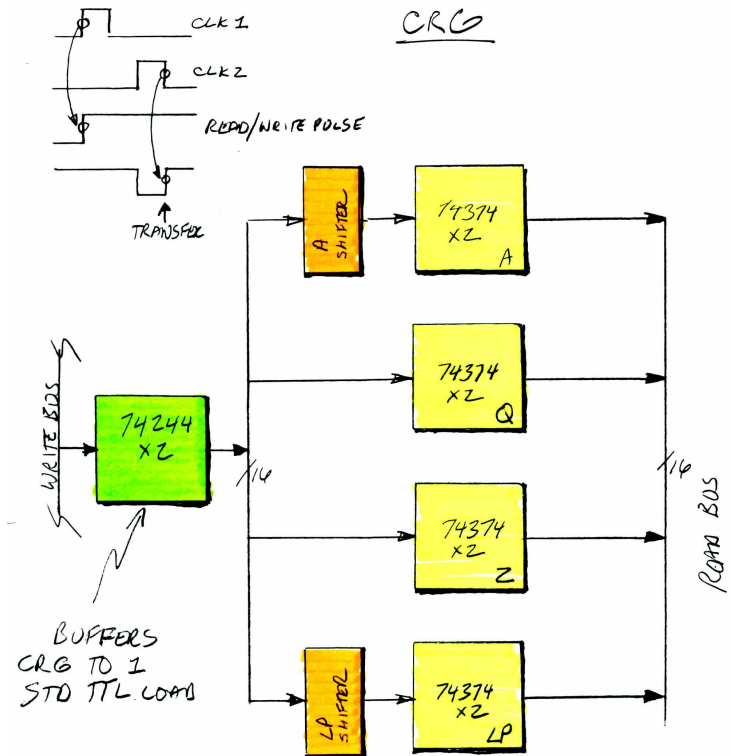




# CRG (Central Register)

The AGC has four 16-bit registers for general computational use. These are called the "central registers":

- A: the 16-bit accumulator, used for general computation.
- Z: the 16-bit program counter, which contains the address of the next instruction to be executed.
- Q: the 16-bit register used to hold the remainder in the DV instruction, and to hold the return address after TC instructions.
- LP: the 16-bit register used to hold the lower product after MP instructions.



## Register A and LP shifters

In addition to "normal" control pulses that write each line of the write bus into the corresponding bit of the registers, the A and LP registers have special write control pulses that shift bits:

The WALP control pulse bit-shifts into the A and LP register. The table below shows how the shifter works. The row of 16 comma-separated entries represent bits in the register. The leftmost position is the register MSB, rightmost position is the LSB. The entry shows the bit of the WRITE bus that's mapped onto that register bit by the shifter.

For the WALP pulse, bit 1 of the WRITE bus (B1) is written into bit 14 of the LP register. "BX" means leave that bit of the register alone (don't change it).

The same WALP pulse causes bit 2 of the WRITE bus (B2) to be written into the lowest bit of the A register, bit 3 of the WRITE bus (B3) to be written into the next bit, and so forth. "US" (uncorrected sign) is the overflow bit (bit 15) of the WRITE bus. "SG" is the 1's complement sign (bit 16) from the WRITE bus.

WALP for register LP:

BX, BX, B1, BX, BX, BX, BX, BX, BX, BX, BX, BX, BX, BX, BX, BX

WALP for register A:

SG, SG, US, B14, B13, B12, B11, B10, B9, B8, B7, B6, B5, B4, B3, B2

Similarly, the WLP control pulse bit-shifts the WRITE bus into the LP register as follows (D0 on bit 14 of the LP register means force the bit to zero):

WLP

B1, B1, D0, B14, B13, B12, B11, B10, B9, B8, B7, B6, B5, B4, B3, B2

The logic design for handling bit 14 of the LP register, which takes control inputs from WALP, WLP, and WA3 is shown here. WA3 is identical to WLP.

Depending upon the WALP, WLP, or WA3 inputs, bit 14 of LP will either be set to B1 of the WRITE bus, or forced to zero.

A	B	C	D	B14	
WA3	WLP	WALP	B1		
0	0	0	0	X	} ILLEGAL
0	0	0	1	X	
0	0	1	0	0	
0	0	1	1	0	
0	1	0	0	X	} ILLEGAL
0	1	0	1	X	
0	1	1	0	0	
0	1	1	1	0	
1	0	0	0	X	} ILLEGAL
1	0	0	1	X	
1	0	1	0	0	
1	0	1	1	0	
1	1	0	0	0	
1	1	0	1	1	
1	1	1	0	X	} INACTIVE
1	1	1	1	X	

		CD			
		00	01	11	10
AB	00	X	X		
	01	X	X		
11			1	X	X
10	X	X			

$$B_{14} = C' \cdot D$$

LOGIC DESIGN FOR BIT 14 INPUT TO LP REGISTER

**CRG INPUTS:**

<u>I/F</u>	<u>signal</u>	<u>full name</u>	<u>state definition</u>
<b>CLK:</b>			
	CLK1	CLOCK 1	1=read bus setup; inhibit read bus output data transfer occurs on falling edge
	CLK2	CLOCK 2	
<b>CPM:</b>			
	RA	READ A	0=output A to read bus
	RA0	READ A	0=output A to read bus
	RQ	READ Q	0=output Q to read bus
	RA1	READ Q	0=output Q to read bus
	RZ	READ Z	0=output Z to read bus
	RA2	READ Z	0=output Z to read bus
	RLP	READ LP	0=output LP to read bus
	RA3	READ LP	0=output LP to read bus
	WA	WRITE A	0=load A from write bus
	WA0	WRITE A	0=load A from write bus
	WQ	WRITE Q	0=load Q from write bus
	WA1	WRITE Q	0=load Q from write bus
	WZ	WRITE Z	0=load Z from write bus
	WA2	WRITE Z	0=load Z from write bus
	WALP	WRITE A,LP	0=load A,LP from write bus
	WLP	WRITE LP	0=load LP from write bus
	GENRST	GENERAL RESET	0=General Reset

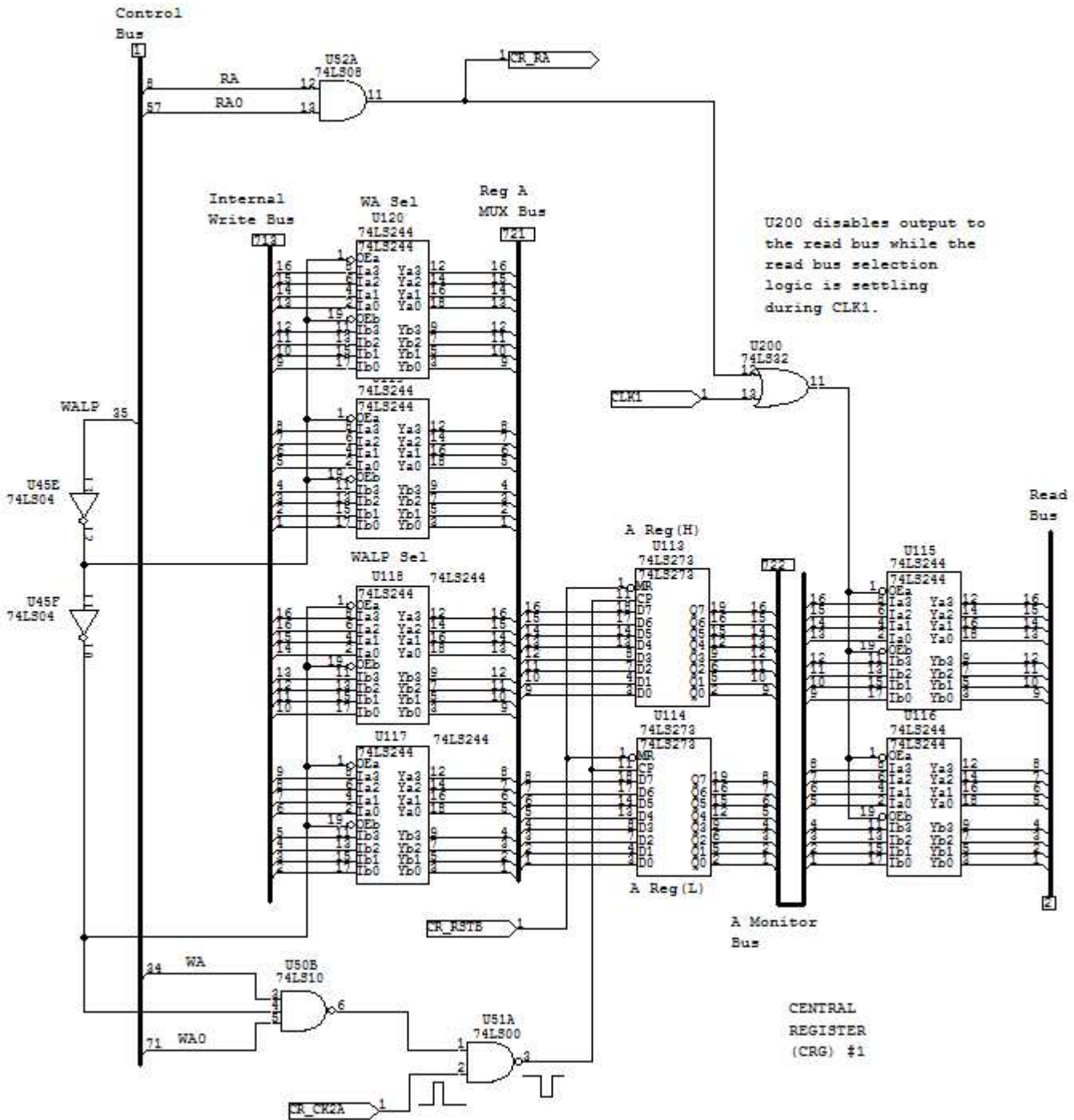


**WBUS:**

WB_01	WRITE BUS 01	
...		
WB_14	WRITE BUS 14	
WB_15	WRITE BUS 15	US (overflow) bit for write bus
WB_16	WRITE BUS 16	SG (sign) bit for write bus

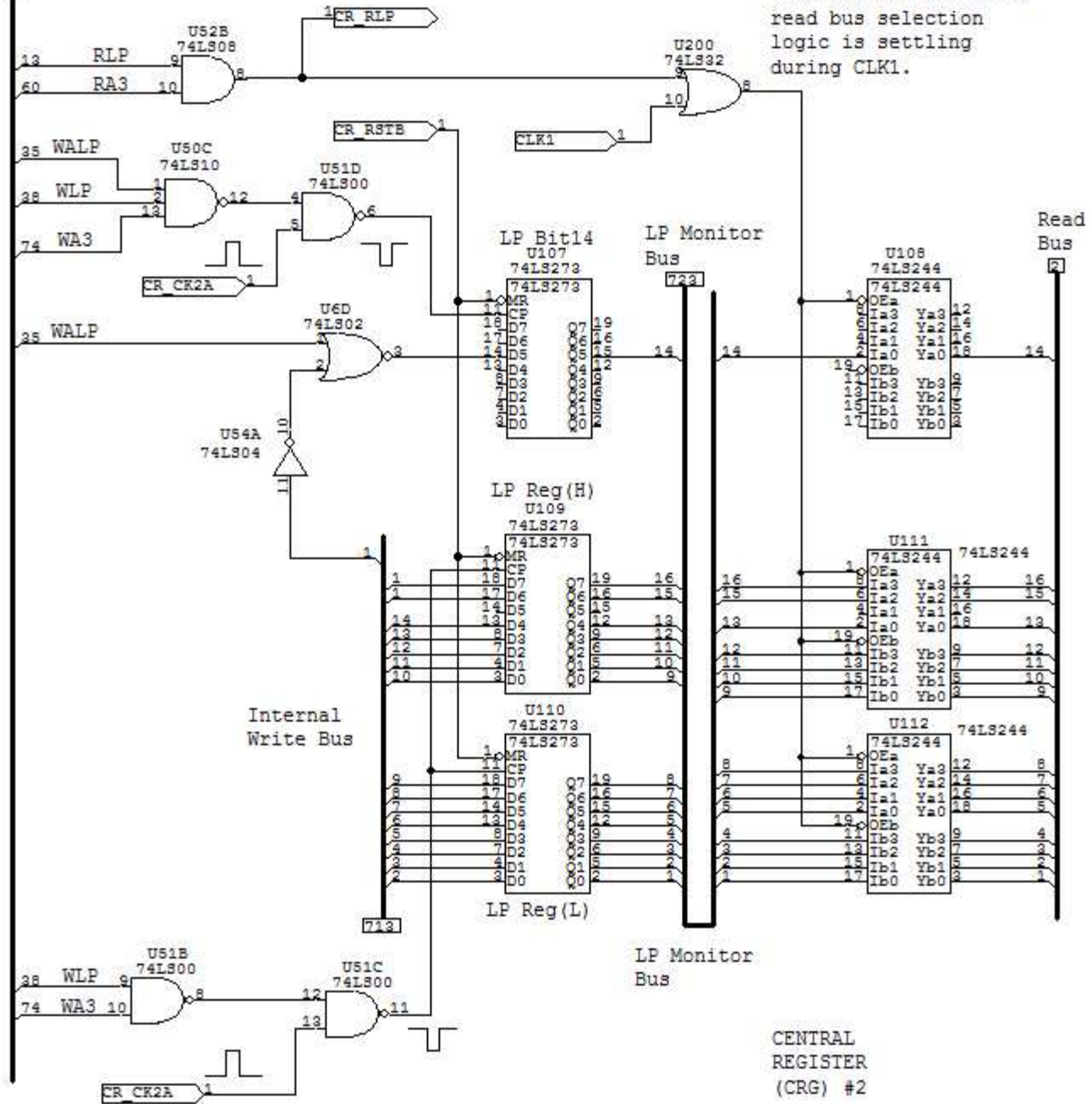
**MBF OUTPUTS:**

<u>I/F</u>	<u>signal</u>	<u>full name</u>	<u>state definition</u>
<b>RBUS:</b>			
	RB_01	READ BUS 01	
	...		
	RB_14	READ BUS 14	
	RB_15	READ BUS 15	US (overflow) bit for read/write bus
	RB_16	READ BUS 16	SG (sign) bit for read/write bus
	BUSY	READ BUS BUSY	0=output enabled to read bus



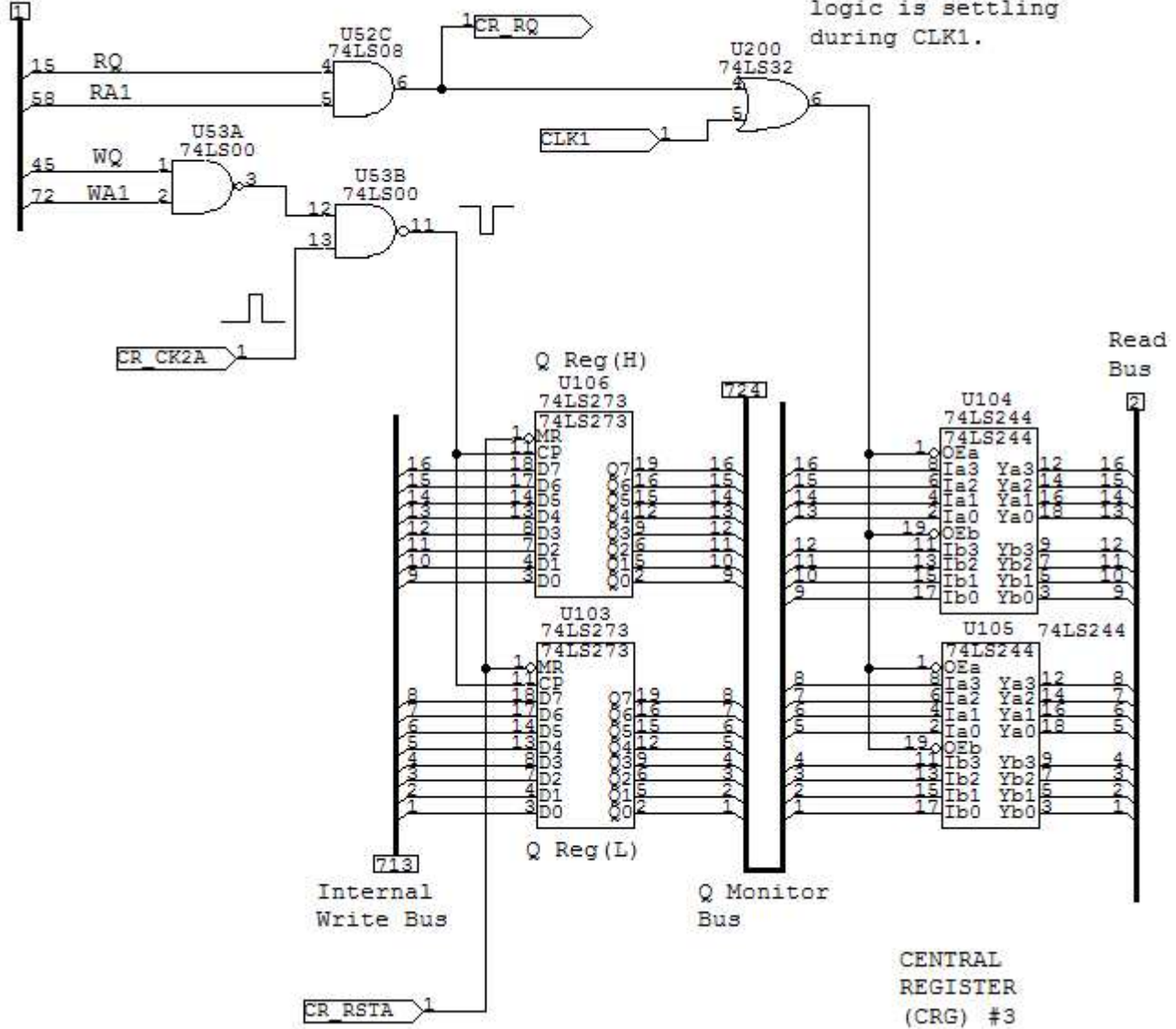
Control Bus

U200 disables output to the read bus while the read bus selection logic is settling during CLK1.



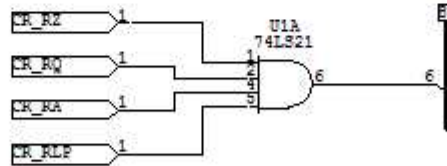
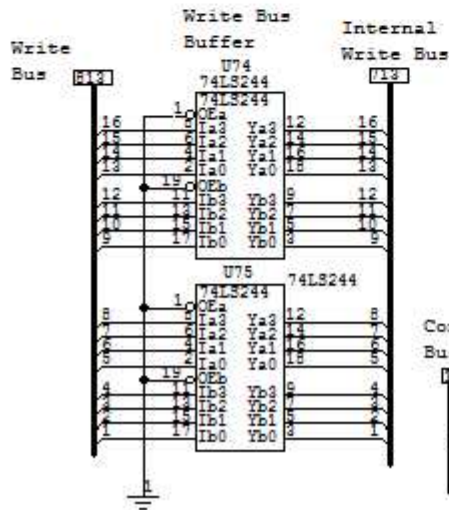
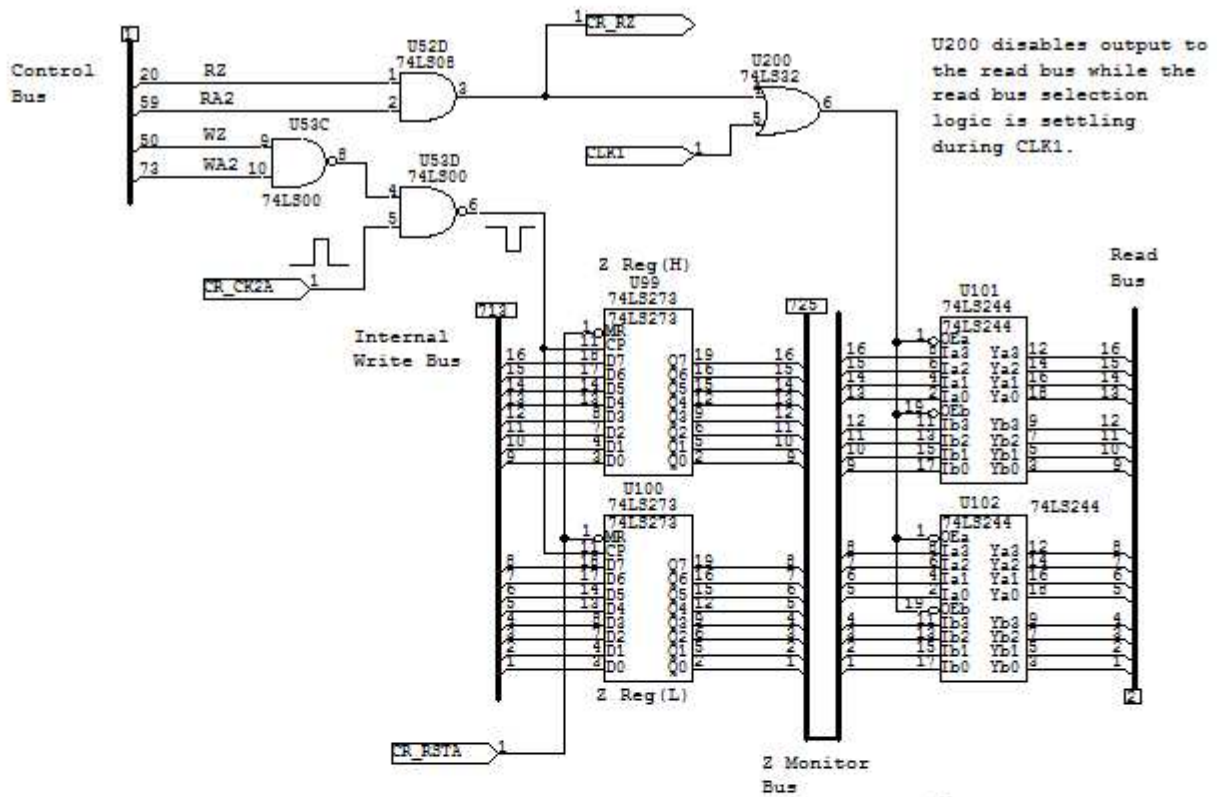
Control Bus

U200 disables output to the read bus while the read bus selection logic is settling during CLK1.

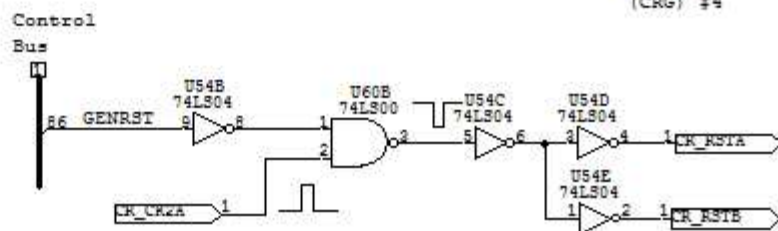


CENTRAL REGISTER (CRG) #3





CENTRAL REGISTER (CRG) #4



## INT (Interrupt Priority)

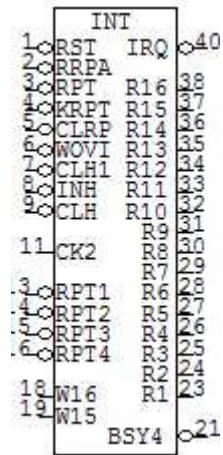
The original AGC had five vectored interrupts. This recreation implements the following 3:

- RUPT1        Also called T3RUPT because it's triggered by overflow of the TIME3 priority counter.
  
- RUPT3        Also called T4RUPT because it's triggered by overflow of the TIME4 priority counter. Because the interrupt is used by software to update the DSKY display at regular intervals, it's sometimes called DSRUPT.
  
- RUPT4        Triggered by a key press from the user's keyboard. Also called KEYRUPT.

The AGC software responds to each interrupt by temporarily suspending the current program, executing a short interrupt service routine, and then resuming the interrupted program.

## INT INPUTS:

<u>I/F</u>	<u>signal</u>	<u>full name</u>	<u>state definition</u>
CLK:	CLK1	CLOCK 1	1=read bus setup; inhibit read bus out data transfer occurs on falling edge
	CLK2	CLOCK 2	
CPM:	GENRST	GENERAL RESET	0=reset INT registers
	RRPA	READ RUPT ADDRESS	0=output RPCELL address (2004,2010,2014,2020,2024) to read bus
	RPT	READ RUPT OPCODE	0=load RUPT opcode into RPCELL register
	KRPT	KNOCK DOWN RUPT PRIO	0=reset RUPT latch currently selected by RPCELL register
	CLRP	CLEAR RPCELL	0=clear RPCELL register
	WOVI	WRITE OVF RUPT INH	0=test overflow; if overflow, inhibit interrupt (set INHINT1)
	CLINH1	CLEAR INHINT1	0=clear INHINT1 register
	INH	SET INHINT	0=set INHINT register
	CLINH	CLEAR INHINT	0=clear INHINT register
WBUS:	WB_15	WRITE BUS 15	US (overflow) bit for write bus
	WB_16	WRITE BUS 16	SG (sign) bit for write bus
CTR/KBD:	RUPT1	INTERRUPT 1	0=trigger interrupt 1 (2004 octal; TIME3 overflow)
	RUPT3	INTERRUPT 3	0=trigger interrupt 3 (2014 octal; TIME4 overflow)
	RUPT4	INTERRUPT 4	0=trigger interrupt 4 (negative edge) (2020 octal; keyboard activity)



Note: interrupt cells above RUPT4 not implemented.

## INT OUTPUTS:

<u>I/F</u>	<u>signal</u>	<u>full name</u>	<u>state definition</u>
SEQ:	IRQ	INT RQST	0=interrupt requested. Active if allof the following are true: a) one or more RUPT FF's are set b) interrupt is not currently being serviced

c) interrupts are not inhibited

RBUS:

RB\_01        READ BUS 01

...

RB\_14        READ BUS 14

RB\_15        READ BUS 15

RB\_16        READ BUS 16

US (overflow) bit for read/write bus

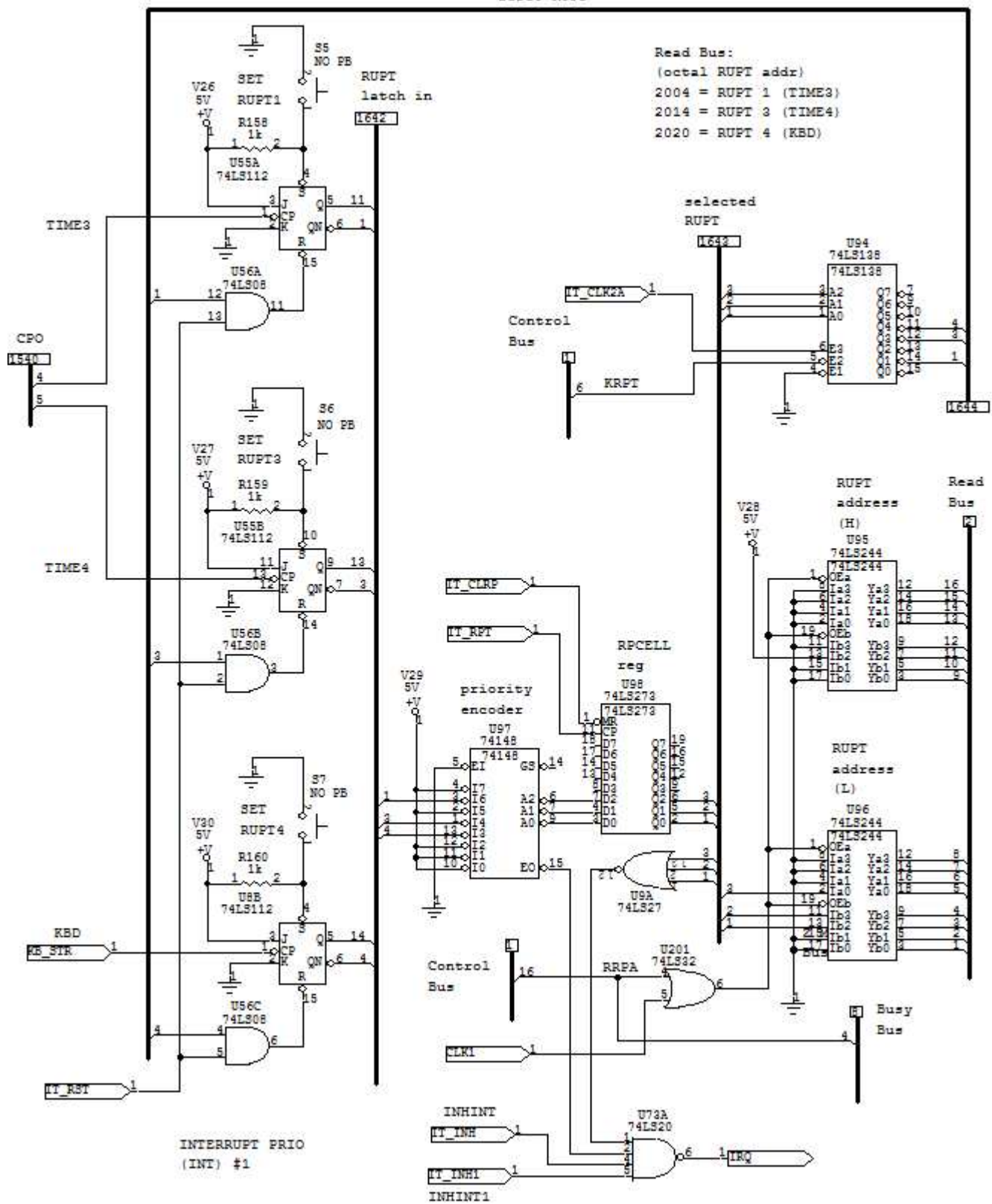
SG (sign) bit for read/write bus

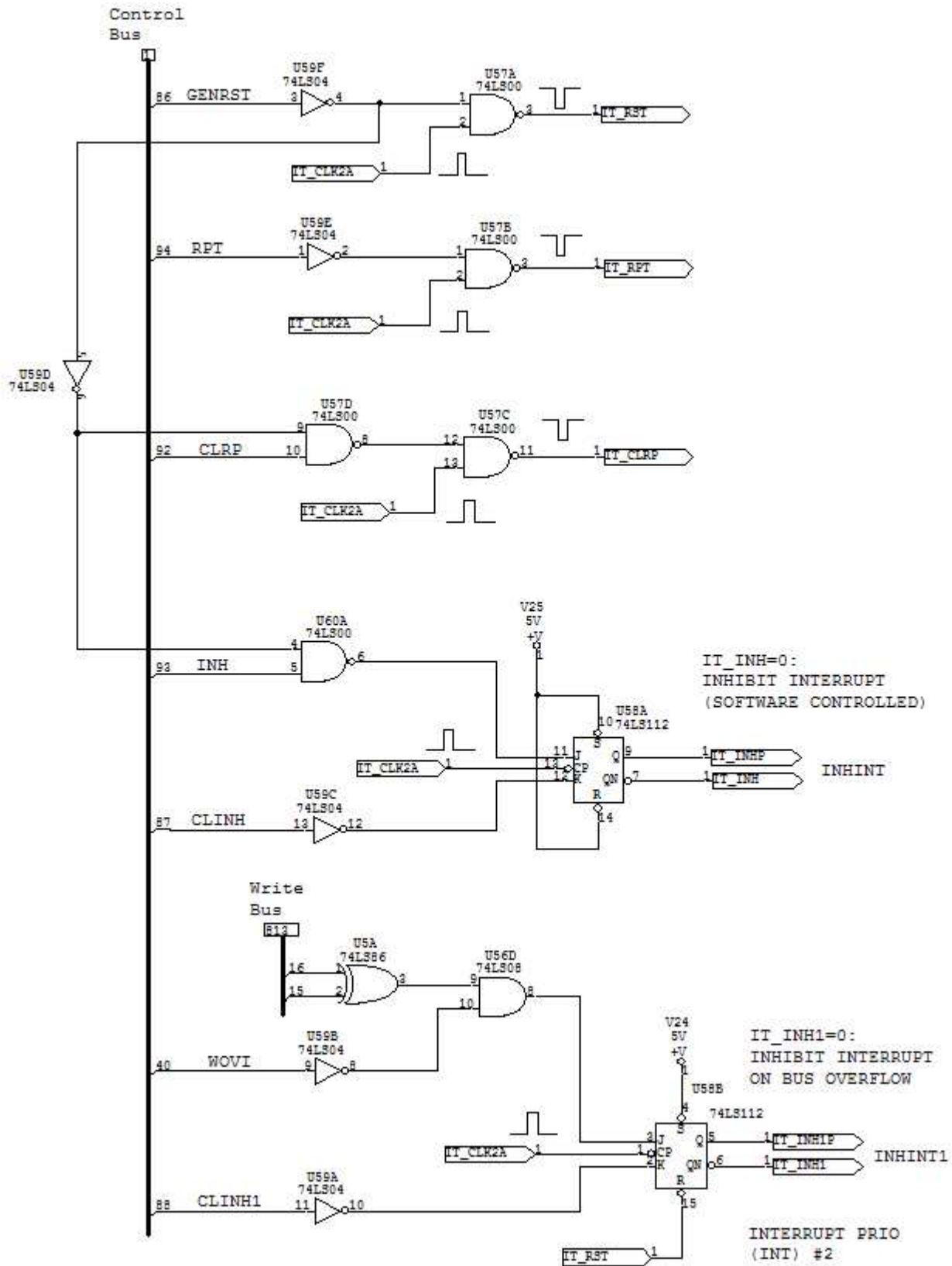
BUSY         READ BUS BUSY

0=output enabled to read bus

reset RUPT

Read Bus:  
(octal RUPT addr)  
2004 = RUPT 1 (TIME3)  
2014 = RUPT 3 (TIME4)  
2020 = RUPT 4 (KBD)





# CTR (Priority Counter)

The Block I AGC had 20 memory locations dedicated as up/down counters (involuntary counters). The counters would increment or decrement in response to external plus or minus logic signals. Increment (PINC) or decrement (MINC) was handled by one subsequence of microinstructions inserted between any two regular instruction subsequences when counter inputs occurred.

This replica implements 5 counters used by the AGC operating system and user interface:

<u>Counter</u>	<u>Addr</u>	<u>Description</u>
OVCTR	34	An overflow counter incremented (PINC) or decremented (MINC) when overflow conditions occur during certain instructions.
TIME2	35	The high-order bits of the AGC clock; incremented (PINC) by overflow of TIME1.
TIME1	36	The low-order bits of the AGC clock; incremented (PINC) by a 100Hz signal from the SCALER (SCL) in the control module (CTL).
TIME3	37	A general purpose timer incremented by a 100Hz signal from the SCALER (SCL) in the control module (CTL).
TIME4	40	A special purpose timer used for software update of the DSKY display. Incremented by a 100Hz signal from the SCALER (SCL) in the control module (CTL).

## INT/CTR interface

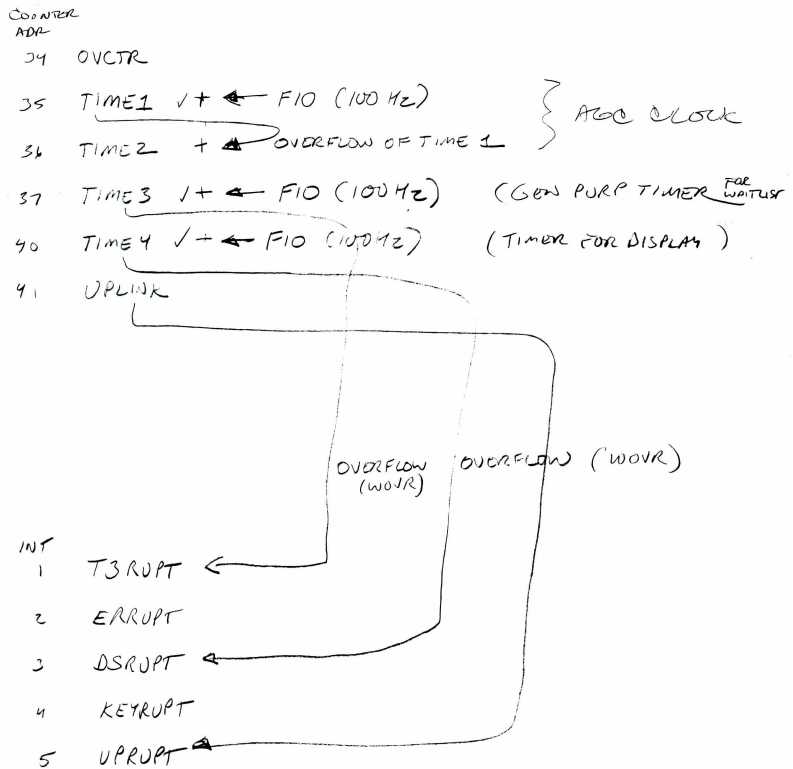
This chart shows how counter overflows are handled. F10 (from the scaler) increments TIME1, TIME3, and TIME4.

A positive overflow of TIME1 causes an increment of TIME2.

Positive overflow of TIME3 triggers a T3RUPT interrupt.

Positive overflow of TIME4 triggers a DSRUPT (T4RUPT) interrupt.

The addresses of TIME1 and TIME2 are reversed for Block II. This chart shows the Block I order, but my replica used the Block II order for compatibility with the COLOSSUS flight software.



## CTR Design Problem

During unit testing, I uncovered a bug in my implementation of CTR: The plus inputs set the P-cell which eventually triggers a PINC subsequence and increments the counter. Minus inputs set the M-cell, which triggers a MINC subsequence and decrements the counter. Near-simultaneous plus and minus inputs should cancel out, producing no change in count.

As a consequence, the counter logic was designed so that, if the P- and M-cells are both set, no counter sequence (PINC or MINC) is selected. The problem is, the P- and M-cells are reset by WOV, which only occurs in PINC or MINC. So, if both cells are set--and therefore, no subsequence is selected--WOV is never issued, the cells never reset, and all counting activity is disabled for that counter.

Here are the options I developed fixing the design:

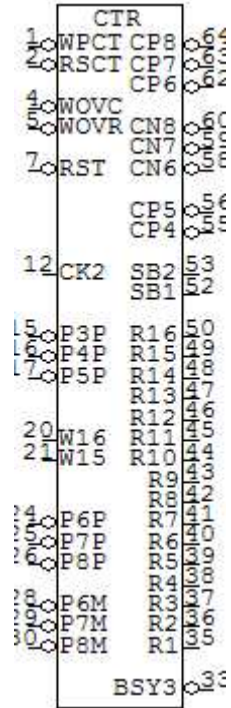
- a) Create a new subsequence similar to PINC or MINC that issues WOV, but does not change the counter. Select this new subsequence when the P- and M-cells are simultaneously set.
- b) Allow the P- and M-cells to trigger PINC and MINC sequences, resulting in a net change of zero to the counter.
- c) Leave the design as-is, because the M-cell is only used for OVCTR in my implementation, and therefore, the problem can never occur.

I choose option C to avoid redesign and re-unit-testing of CTR.



**CTR INPUTS:**

<u>I/E</u>	<u>signal</u>	<u>full name</u>	<u>state definition</u>
CLK:	CLK1	CLOCK 1	1=read bus setup; inhibit read bus out data transfer occurs on falling edge
	CLK2	CLOCK 2	
CPM:	GENRST	GENERAL RESET	0=reset CTR registers
	WPCTR	WRITE PSEQ	0=write sequence into PSEQ
	RSCT	READ PCELL ADDRESS	0=output PCELL address (034-043) to read bus
	WOVC	WRITE OVRFLOW CNTR	0=test overflow and inc/dec OVCTR
	WOVR	WRITE OVERFLOW	0=clear selected PCELL and handle counter overflow (if any)
WBUS:	WB_15	WRITE BUS 15	US (overflow) bit for write bus
	WB_16	WRITE BUS 16	SG (sign) bit for write bus
EXTERNAL:	P3P	P3 CELL + COUNT	0=count up P3 counter (036 octal; TIME1)
	P4P	P4 CELL + COUNT	0=count up P4 counter (037 octal; TIME3)
	P5P	P5 CELL + COUNT	0=count up P5 counter (040 octal, TIME4)



Note: priority cells 6-20 not implemented.

**CTR OUTPUTS:**

<u>I/E</u>	<u>signal</u>	<u>full name</u>	<u>state definition</u>
COUNTER OVERFLOW:			
	CPO_04	P4 + OVERFLOW	0=P4 cell pos ovf (during WOVR) note: TIME3 + overflow; connect to INT subsystem to trigger T3RUPT interrupt.
	CPO_05	P5 + OVERFLOW	0=P5 cell pos ovf (during WOVR)

note: TIME4 + overflow; connect to INT subsystem to trigger T4RUPT (DSRUPT) interrupt.

SEQ:

SB_01	SUB SEL 01	SB_01 is LSB; SB_02 is MSB
SB_02	SUB SEL 02	00=no counter; 01=PINC; 10=MINC

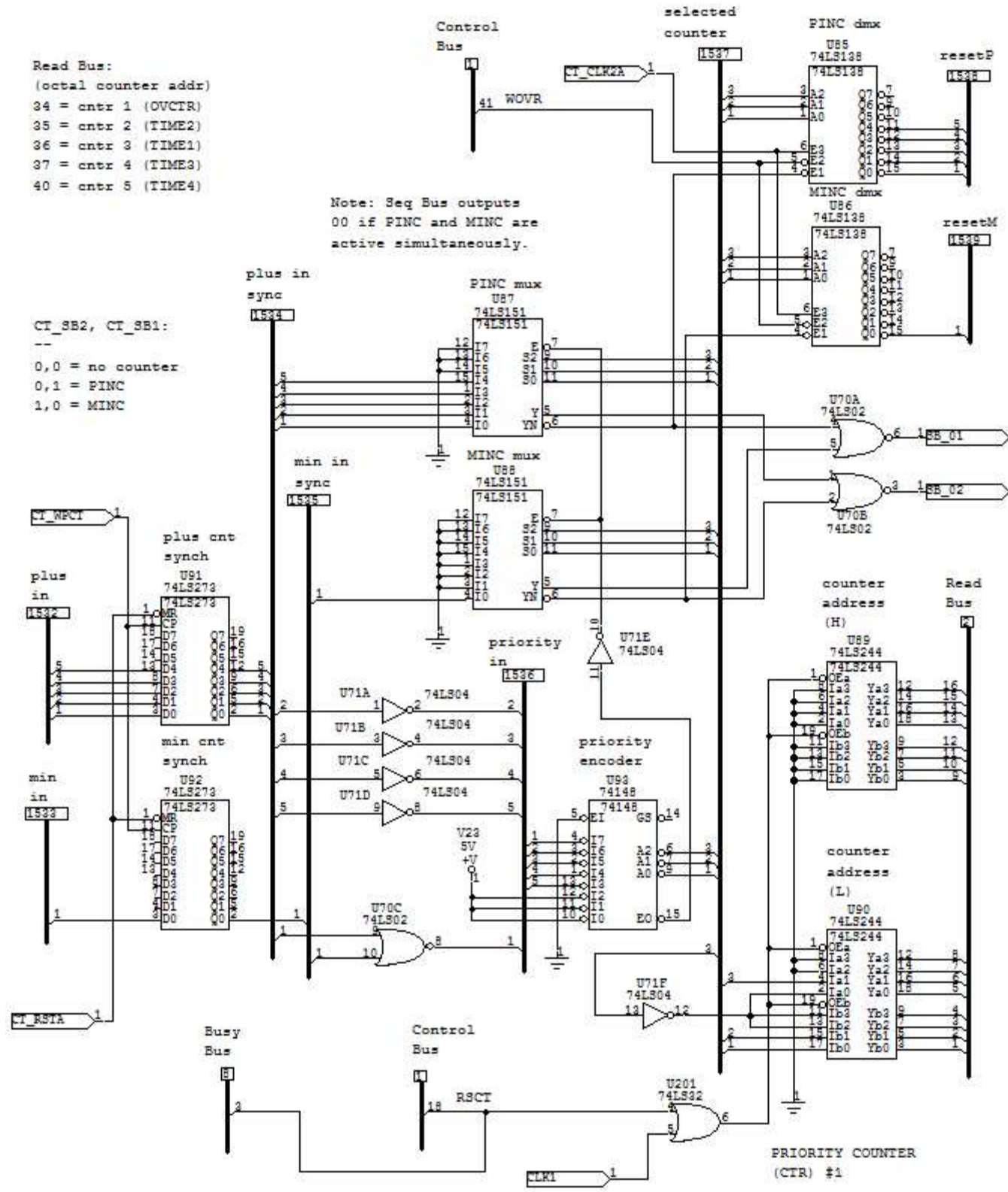
RBUS:

RB_01	READ BUS 01	
...		
RB_14	READ BUS 14	
RB_15	READ BUS 15	US (overflow) bit for read/write bus
RB_16	READ BUS 16	SG (sign) bit for read/write bus
BUSY	READ BUS BUSY	0=output enabled to read bus

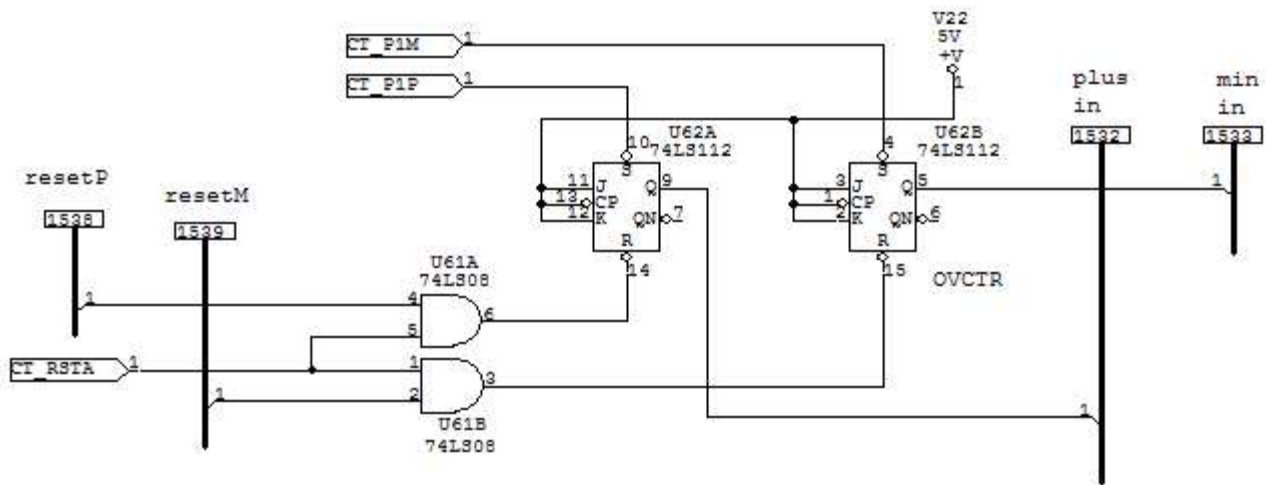
Read Bus:  
 (octal counter addr)  
 34 = cntr 1 (OVCTR)  
 35 = cntr 2 (TIME2)  
 36 = cntr 3 (TIME1)  
 37 = cntr 4 (TIME3)  
 40 = cntr 5 (TIME4)

CT\_SB2, CT\_SB1:  
 --  
 0,0 = no counter  
 0,1 = PINC  
 1,0 = MINC

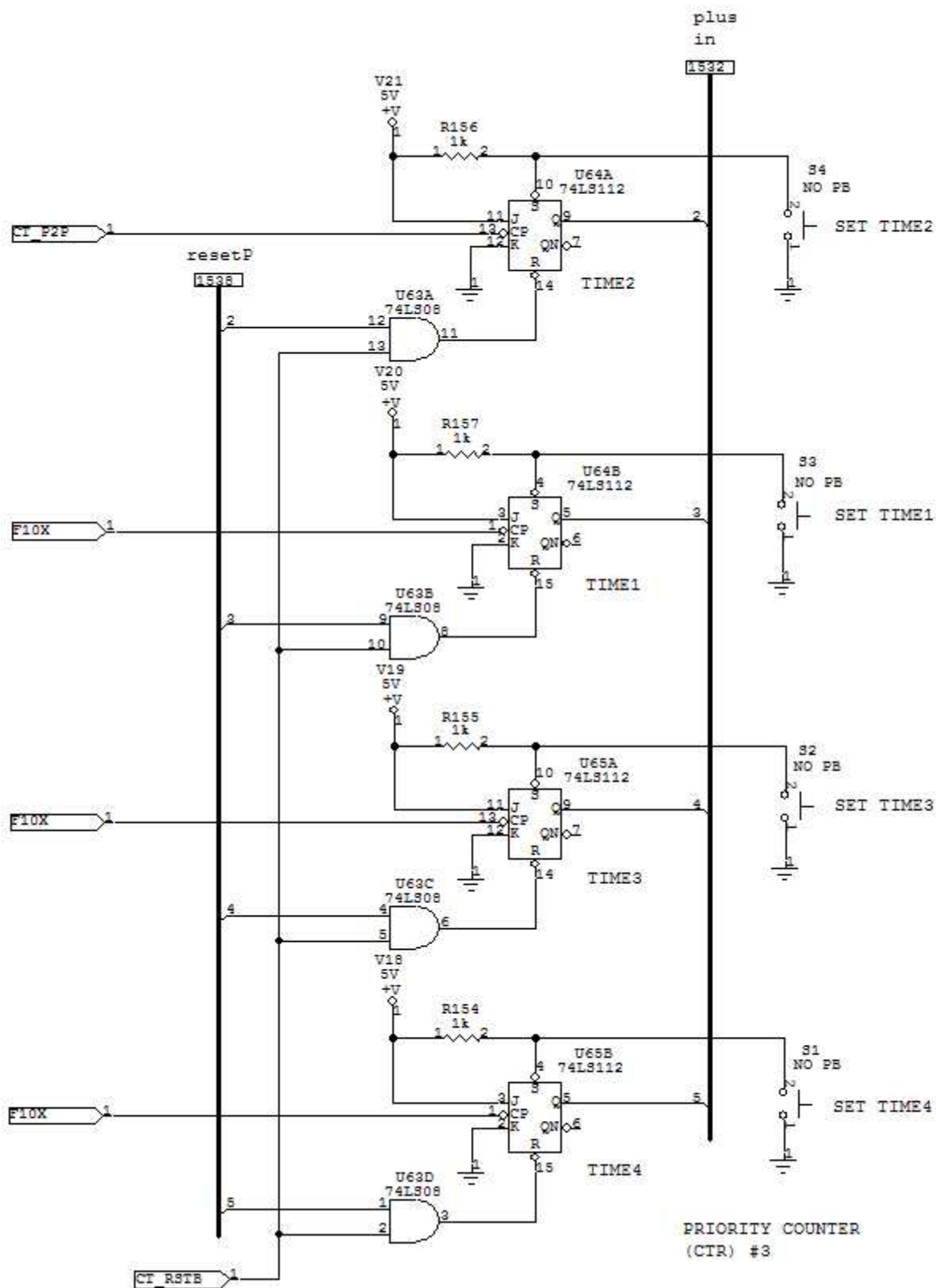
Note: Seq Bus outputs  
 00 if PINC and MINC are  
 active simultaneously.

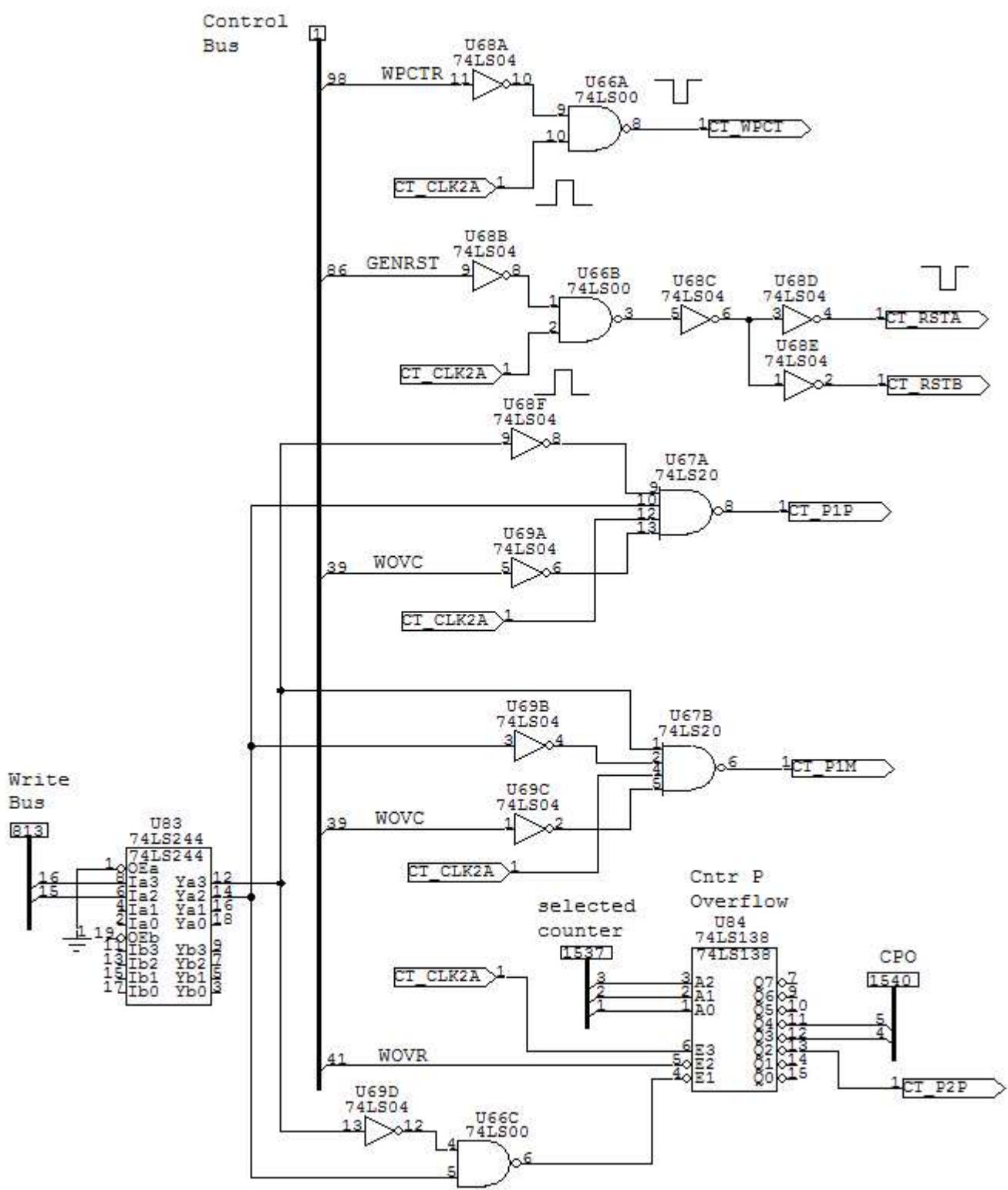


PRIORITY COUNTER  
 (CTR) #1



PRIORITY COUNTER  
(CTR) #2





PRIORITY COUNTER  
(CTR) #4

# Fabrication

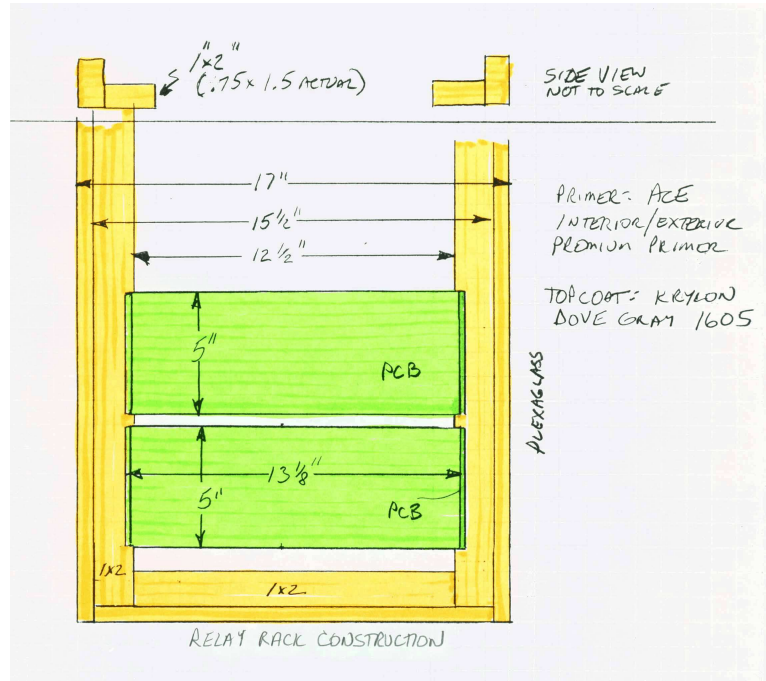
The PROC module is (4) 13"x5" circuit boards, and 1 control panel.

## Module Rack

The module framework is designed to resemble a relay rack, but scaled to fit the circuit board dimensions. It is constructed out of 1"x2" pine and spray-painted semi-gloss gray.

Circuit boards are mounted to the rack by 2 phillips screws at either end. Nylon spacers (1/4") are used as standoffs to hold the board edges above the rack. The boards are mounted so the chips are in the back and the pins wiring are visible from the front.

Power is distributed by 2 heavy aluminum bus bars mounted vertically, one per side, on the back of the module. Machine screws are mounted through the bus bars at evenly-spaced intervals to provide connection points for the boards.

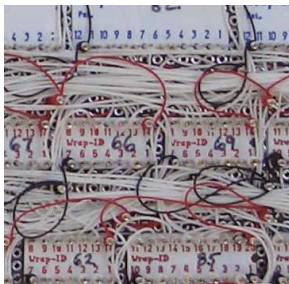


Solid copper wire (24 gauge) connects the boards to the bus bars. Ring terminals are used on the bus bar side of the connection. On the circuit board side, the wires are soldered directly to the supply rails.

Materials were purchased from Home Depot, ACE Hardware, and Radio Shack.

## Circuit Boards

The circuit boards are 13"x5" general purpose prototyping boards, epoxy glass with double-side plated through pads on 0.1" centers (JAMECO 21477CL).



ICs are mounted in level 3 machine tooled wire-wrap sockets: 8, 14, 16, 20, 24, and 28 pin (JAMECO). Each socket has the pin-out labeled with a wire-wrap socket ID marker, which slips onto the socket before wrapping (JAMECO). The part number is written onto the ID marker.

Sockets are arranged in 4 horizontal rows on each board, with about 10 sockets per row.

Power is distributed on the back-side of each board by bare 24-gauge solid copper wire supply rails soldered at equal intervals to Klipwrap terminals: 3-prong terminals with a square tail for wire-wrapping (JAMECO 34163CL). A +5V rail runs above each row of sockets and a ground rail runs below. Each rail connects directly to the aluminum module power bus using a ring tail connector.

On the pin side of the board, all connections are made with 30 AWG Kynar wire-wrap wire (JAMECO). Red wire is used for direct connections to the +5V supply rail. Black wire is used for direct connections to ground. White wire is used for everything else.

Power connections from the supply rails to each ICs are double-wrapped. Bypassing capacitors (.1 uf disc ) are soldered across the supply rails at the Klipwrap terminals; about 1 capacitor for every 2 IC packages.

All connections were stripped and hand-wrapped using a Radio Shack hand-wrap tool. As each connection was made, the corresponding line on the schematic was marked with a colored highlighter.

DIP resistor networks (JAMECO) plugged into 20-pin wire-wrap sockets were used as current limiting resistors for the panel indicators.



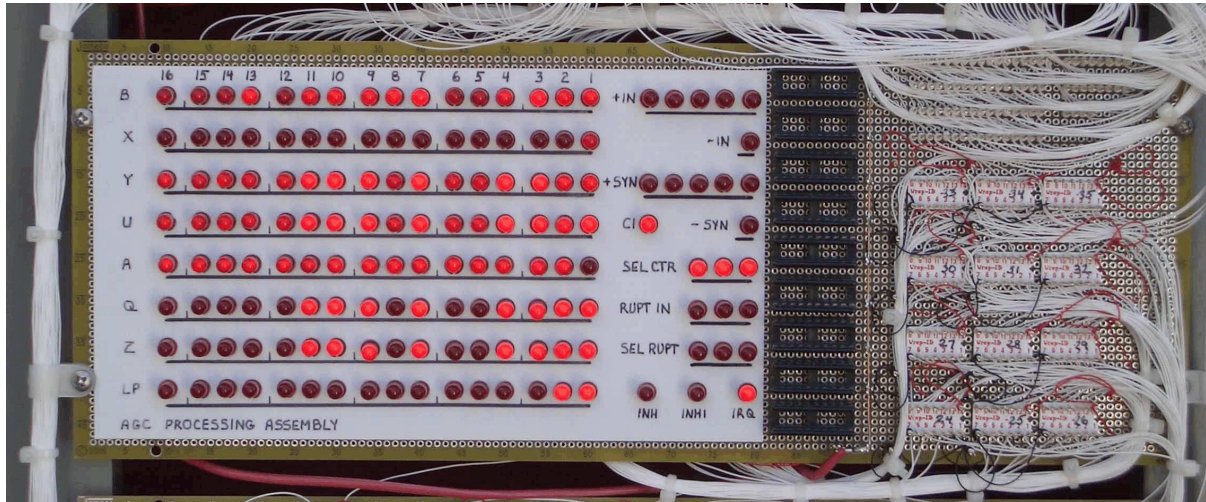
# PROC Printed Circuit Board (PCB) A

The A board contains display drivers for the B board (left side), buffers for the interfaces to external modules (bottom right), and priority counter logic (upper right). Sockets for 3 IDE interface cables to external modules are visible at the bottom.



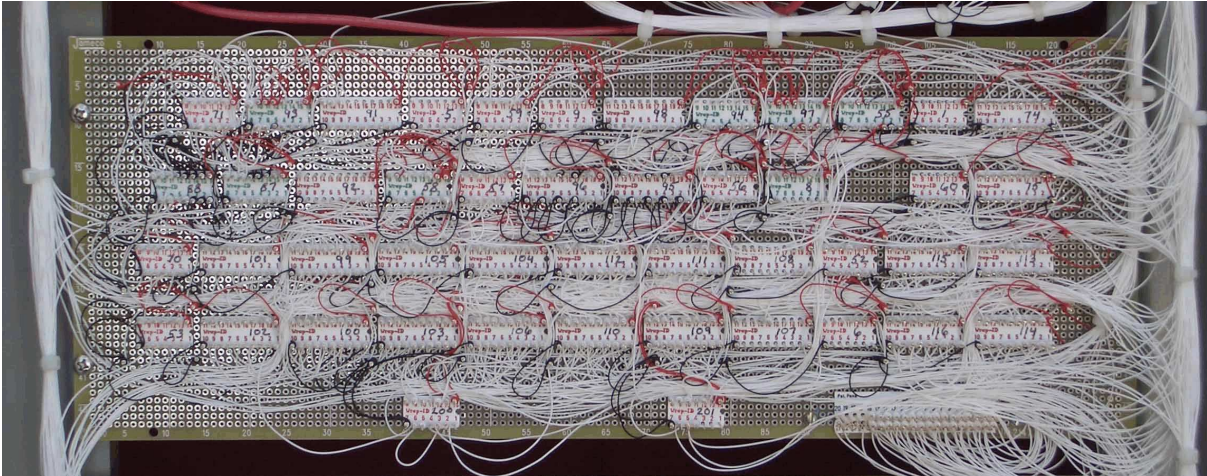
## PROC Printed Circuit Board (PCB) B

The B board contains the display indicators, their current limiting resistor networks, and the open collector drivers. The display panel is a sheet of white styrene plastic. A push pin was used to make holes through the plastic and the LEDs were inserted in rows. The panel was hand-lettered with an indelible marker.



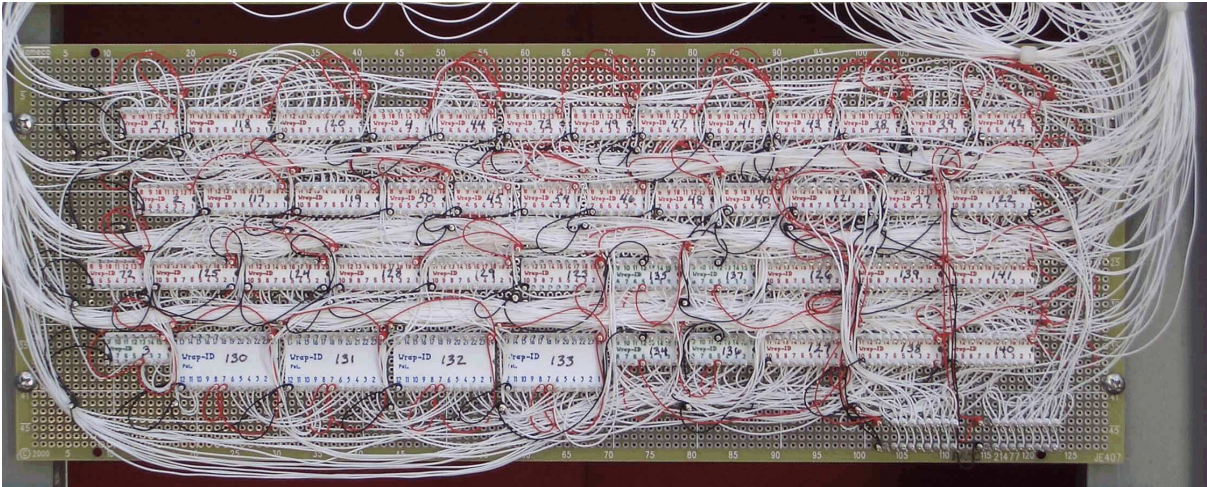
# PROC Printed Circuit Board (PCB) C

The C board contains the the logic for the interrupt (INT) subsystem (upper half of the board), and the central registers (CRG; lower half of the board).



# PROC Printed Circuit Board (PCB) D

The D board contains the ALU logic. The large 74181 ALU chips are at the bottom right. The four chips that form the ADDER are in the bottom half of the board, slightly to the right of the middle.



## Parts (ICs)

IC's, sockets, PCB's, resistors, capacitors, wire-wrap wire were purchased from JAMECO. IDE wire-wrap sockets were from DigiKey. Wire ties, wire-wrap tools, and copper wire were from Radio Shack. IDE ribbon cables were purchased from an online computer supplier.

74LS00	(9)	U60,U66,U57,U53,U51,U44,U46,U47,U42
74LS02	(4)	U70,U6,U49,U2
74LS04	(11)	U54,U45,U37,U68,U7,U36,U69,U71,U59,U48,U72
74LS06	(26)	U28,U29,U30,U31,U27,U26,U23,U24,U25,U22,U21,U20,U18,U19,U17,U16,U15,U12,U13,U14,U11,U10,U35,U34,U33,U32
74LS08	(4)	U63,U61,U56,U52
74LS10	(2)	U50,U43
74LS20	(2)	U67,U73
74LS21	(1)	U1
74LS27	(1)	U9
74LS32	(5)	U4,U41,U40,U39,U38
74LS83	(4)	U134,U135,U136,U137
74LS86	(1)	U5
74LS112	(7)	U65,U64,U62,U58,U55,U8,U3
74LS138	(4)	U84,U85,U86,U94
74LS148	(2)	U93,U97
74LS151	(2)	U87,U88
74LS181	(4)	U130,U131,U132,U133
74LS244	(33)	U74,U75,U76,U77,U78,U79,U80,U81,U82,U83,U89,U90,U95,U96,U101,U102,U104,U105,U108,U111,U112,U115,U116,U117,U118,U119,U120,U121,U122,U124,U125,U128,U129
74LS273	(19)	U91,U92,U98,U99,U100,U103,U106,U107,U109,U110,U113,U114,U123,U126,U127,U138,U139,U140,U141

## Power Budget

	<u>qty</u>	<u>mA (ea)</u>	<u>mA (tot)</u>
74LS00	9	2.4	21.6
74LS02	4	2.4	9.6
74LS04	11	3.6	39.6
74LS06	26	3.6	93.6
74LS08	4	4.4	17.6
74LS10	2	1.8	3.6
74LS20	2	1.2	2.4
74LS21	1	2.2	2.2
74LS27	1	3.4	3.4
74LS32	5	4.9	24.5
74LS83	4	22.0	88.0
74LS86	1	6.1	6.1
74LS112	7	4.0	28.0
74LS138	4	6.3	25.2
74LS148	2	12.0	24.0
74LS151	2	6.0	12.0
74LS181	4	21.0	84.0
74LS244	33	32.0	1056.0

74LS273	19	17.0	323.0
LED	153	20.0	3060.0
			-----
			4.9 Amps total
			1.9 Amps (excluding LEDs)